

**Secure Implementation Experiments:
Do Strategy-proof Mechanisms Really Work?†**

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ABSTRACT

Strategy-proofness, requiring that truth-telling is a dominant strategy, is a standard concept used in social choice theory. Saijo et al. (2003) argue that this concept has serious drawbacks. In particular, announcing one's true preference may not be a unique dominant strategy, and almost all strategy-proof mechanisms have a continuum of Nash equilibria. For only a subset of strategy-proof mechanisms do the set of Nash equilibria and the set of dominant strategy equilibria coincide. For example, this double coincidence occurs in the Groves mechanism when preferences are single-peaked. We report experiments using two strategy-proof mechanisms where one of them has a large number of Nash equilibria, but the other has a unique Nash equilibrium. We found clear differences in the rate of dominant strategy play between the two. *Journal of Economic Literature* Classification Number: C92, D71, D78, and H41.

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1. Introduction

Strategy-proofness, requiring that truth-telling is a dominant strategy, is a standard concept that has been used in the design of a variety of mechanisms for social choice as well as for eliciting values for non-market goods. Its main appeal is that it relies on what would seem to be one of the most basic game-theoretic notions and apparently innocuous assumptions for behavior: that players adopt dominant strategies. Theorists often fail to recognize, however, that laboratory evidence calls into question the descriptive relevance of this assumption. For example, Attiyeh, Franciosi, and Isaac (2000) and Kawagoe and Mori (2001) report pivotal mechanism experiments in which subjects adopt dominant strategies less than half the time, and Kagel, Harstad, and Levin (1987), Kagel and Levin (1993) and Harstad (2000) report second price auction experiments in which most bids do not reveal true value. Attiyeh, Franciosi, and Isaac (2000) conclude pessimistically (p. 112) “we do not believe that the pivot mechanism warrants further practical consideration... This is due to the fundamental failure of the mechanism, in our laboratory experiments, to induce truthful value revelation”.

Experimentalists sometimes argue that players who use weakly dominated strategies must suffer from confusion due to the complexity of the mechanism and the non-transparency of the dominant strategy. But in fact, neither “epistemic” (deductive) nor “evolutionary” (dynamic) models provide unambiguous support for the elimination of weakly dominated strategies. If each player is perfectly rational and can deduce what strategies the opponent will use, then the outcome of the game must be a Nash equilibrium (Aumann and Brandenburger (1995)), but there is nothing that forces a player to eliminate weakly dominated strategies. In a dynamic analysis, the behavior of boundedly rational players is changing over time. While the rest points of dynamic processes such as fictitious play must be Nash equilibria, there is no guarantee that weakly dominated strategies will be eliminated. Intuitively, the feedback the players receive

may be very weak because the use of a weakly dominated strategy may not cause any loss in payoff. Binmore, Gale and Samuelson (1995) and Kagel and Levin (1993) argue that this weak feedback effect can explain some experimental results, and Cabrales and Ponti (2000) discuss the implications for mechanism design. Of course, epistemic and evolutive models do provide clear-cut support for the elimination of *strictly* dominated strategies. The problem is that very few social choice rules are implementable in strictly dominant strategies.

Motivated by this problem, Saijo, Sjöström and Yamato (2003) developed a new concept called *secure* implementation. A social choice function is securely implementable if there exists a mechanism (game form) that implements it in dominant strategy equilibria, and the set of dominant strategy equilibrium outcomes and the set of Nash equilibrium outcomes coincide. That is, all Nash equilibrium outcomes must be socially optimal in a secure mechanism. The current paper takes a first step towards establishing the empirical significance of these ideas. We report a new experiment comparing the rate of dominant strategy adoption for the pivotal mechanism (where implementation is not secure) and for the Groves-Clarke mechanism when preferences are single-peaked (where implementation is secure). Our results indicate that subjects play dominant strategies significantly more often in the secure Groves-Clarke mechanism than in the non-secure pivotal mechanism, even though we have simplified both mechanisms with context-free payoff tables. Our findings suggest that the highly pessimistic conclusion of Attiyeh, Franciosi, and Isaac (2000) should be modified to allow the possibility that a Groves-Clarke mechanism can perform satisfactorily in environments where implementation is secure.

The practical relevance of mechanism design will increase as more mechanisms are implemented in the field. Auctions provide an important example. The English (ascending price) auction is a secure mechanism that has been used since at least 500 B.C. in Babylon

(Cassady, 1967). Theorists have noted the strategic equivalence between English and second price auctions since Vickrey (1961), but for some information conditions the second price auction is strategy-proof but not securely implementable. Until recently the second price auction has not been adopted in the field, although this is likely to change as online auctions grow in importance. Bidders in online auctions at eBay and Amazon can submit a reservation price (called a proxy bid) early in the auction, and if this bid is highest then this bidder wins the auction and pays only the minimum bid increment above the second-highest submitted price. This institution shares a number of incentive features of theoretical second price auctions, although as currently implemented submitting one's reservation price is generally not a dominant strategy (Roth and Ockenfels, 2002).

But the adoption of true sealed-bid second price auctions may grow over time, particularly for intermediate goods and in procurement ("business-to-business") transactions.¹ But as we illustrate in Section 3, for some information conditions the second price auction has "bad" Nash equilibrium outcomes that are Pareto-inferior to the dominant strategy equilibrium outcome. This suggests that proponents of second price auctions may want to be more cautious when proposing them for online markets or to elicit valuations for non-market goods.

The remainder of the paper is organized as follows. Section 2 presents a brief review of the laboratory evidence on strategy-proof mechanisms. Section 3 gives examples of two well-known strategy-proof mechanisms that have a continuum of Nash equilibria, including equilibria other than the dominant strategy equilibrium that theorists usually focus on. We characterize secure implementability in Section 4 for the case of two agents and quasi-linear preferences that is relevant for our experiment (Saijo et al. (2003) presents results for more

¹ Some examples where sealed-bid second price auctions have been introduced for business-to-consumer and consumer-to-consumer transactions include qconlineauction.com and grab-a-deal.com.

general conditions). Section 5 describes the experimental environment and Section 6 contains the experimental results. Section 7 provides concluding remarks.

2. Experimental Results on Strategy-Proof Mechanisms

Until recently, most of the experimental studies of strategy-proof mechanisms have considered the second price auction (Vickrey, 1961). For example, Coppinger, Smith and Titus (1980) studied the relationship between Dutch, English, first price sealed-bid and second price sealed-bid auctions. Bidders in both the English and the second price auction have a dominant strategy to fully reveal their resale value in their bid (or reveal their value in their “drop-out price” in the case of the English auction). Bidders in Coppinger et al.’s (oral) English auctions typically dropped out of the bidding when predicted, so prices corresponded to the equilibrium prediction – the second-highest bidders’ resale value. Similarly, Kagel, Harstad and Levin (1987) show that bidders in English (clock) auctions lock on to the dominant strategy of bidding equal to value after a few periods of initially overbidding.

Bidders in Coppinger et al.’s second price auctions did not bid above their resale value, but this is clearly because of the artificial prohibition of bids above resale value imposed in these initial experiments. Kagel and Levin (1993) find that 58 to 67 percent of second price auction bids are greater than resale value, which they attribute to (1) the equilibrium bidding strategy being less transparent than in the English auction and (2) learning feedback to discourage overbidding is weak under sealed-bid procedures because typically the overbidding is not “punished” with losses. Harstad (2000) also documents rather severe overbidding in second price auctions that does not decline over time but that may be less pronounced when subjects first obtain experience in English auctions. Garratt, Walker and Wooders (2002) show that bidders who are highly experienced in online auctions are no more likely to overbid than to

underbid, but as with inexperienced bidders only very few (roughly 20 percent) of bids are approximately equal to value. Most bids in the Garratt et al. study vary considerably from the bidders' true values, and consequently less than one half of the auctions result in efficient allocations. Overall the data clearly indicate that subjects do not play their dominant strategy, and in all cases the evidence suggests that bidding equal to value is significantly more common in English than in second price auctions.

While the transparency, experience and feedback explanations for the lower frequency of dominant strategy play in the second price auction are all plausible, we propose a complementary explanation. In English auctions with a stage-game structure, the (sub-game perfect) Nash equilibrium outcome coincides with the dominant strategy equilibrium outcome in which bids fully reveal values. But in second price sealed-bid auctions with a one-shot game structure, Nash equilibria that do not coincide with the dominant strategy equilibrium exist and involve overbidding and underbidding. For example, suppose bidder 1 has a value of \$555 and bidder 2 has a value of \$550, and that these values are common knowledge. It is a Nash equilibrium for bidder 1 to bid \$540 and bidder 2 to bid \$560, resulting in the inefficient allocation of the object to bidder 2. Kagel and Levin (1993) and others have noted that overbidding is not discouraged because bidders can bid above values and not lose money. It is precisely this feature of the second price auction institution that causes "bad" Nash equilibria to exist. Saijo et al. (2003) discuss many other examples of strategy-proof mechanisms that also have bad Nash equilibrium outcomes that are Pareto-inferior to the dominant strategy equilibrium outcome.

More recent experiments have studied the pivotal mechanism, which is a strategy-proof

social choice mechanism that is strategically equivalent to the second price auction.^{2,3} In this mechanism an agent pays the amount needed to implement his preferred outcome only if his report is pivotal and changes the chosen outcome. These studies have also documented that subjects frequently do not play dominant strategies. Attiyeh, Franciosi and Isaac (2000) find that less than 10 percent of the bids reveal the subjects' true value for the public good, in a setting where the experimenter explained the mapping of bids to outcomes (and required taxes for the pivotal players) for five- and ten-person groups. Part of the poor performance of this mechanism might be due to subject confusion and the complexity of the pivotal mechanism. Kawagoe and Mori (2001) provide support for this interpretation, using a controlled experiment that manipulates the complexity across treatments. They also find that only a small number of bids (less than 20 percent) reveal true values when the context and complexity of the pivotal mechanism is part of the experiment; but when the mechanism is simplified and represented by (detailed) payoff tables then nearly half of the subjects play the dominant strategy. In the present experiment we also study the pivotal mechanism with detailed payoff tables to help simplify the decision environment and promote equilibrium bids. Although confusion and complexity may be partly responsible for the poor performance of some mechanisms, we will try to go beyond this explanation. We will argue that the existence of multiple Nash equilibria allows us to predict *how* behavior will deviate from the dominant strategy equilibrium. That is, we will identify systematic rather than random deviations from the dominant strategy equilibrium in non-secure mechanisms.

² Another truth-telling mechanism that has been widely employed in experiments is the Becker-DeGroot-Marshak (BDM) mechanism. In this mechanism the subject states a maximum buying price or minimum selling price, but the actual buying or selling price is determined by a randomizing device and the transaction is carried out if it is acceptable giving the subject's reported maximum or minimum. This mechanism is not a game so it is not directly relevant for our study.

³ We do not review here other social choice mechanism experiments like the serial cost sharing mechanism because the researchers have implemented those mechanisms in environments where the Nash equilibria are not in dominant strategies (e.g., Chen, 2003; Dorsey and Razzolini, 1999).

3. Why do Strategy-Proof Mechanisms Not Work Well?

Saijo et al. (2003) show that many of the strategy-proof mechanisms that have been studied in the literature have a continuum of Nash equilibrium outcomes that do not coincide with the dominant strategy equilibrium outcome. In particular, they present a number of examples including the pivotal mechanism for a non-excludable public good, the serial cost sharing mechanism for an excludable public good, the second price auction for an indivisible good, the Condorcet winner voting scheme (a median voter scheme) with single-peaked preferences, and the uniform allocation rule (a fixed-price trading rule) with single-peaked preferences. Besides having a continuum of Nash equilibria, these mechanisms all have bad Nash equilibrium outcomes that are Pareto-inferior to the dominant strategy equilibrium outcome. Here we provide more details for the two examples of such strategy-proof mechanisms that were just summarized from the experimental literature.

Example 1: The pivotal mechanism (Clarke, 1971).

Consider the pivotal mechanism, which is one of the two mechanisms studied in the present experiment, for a two-agent economy with a binary non-excludable public good and quasi-linear preferences. Two agents 1 and 2 are facing a decision whether or not they should produce the public good. Agent i 's true net value of the public good is v_i if it is produced, and her true net value is 0 otherwise ($i = 1, 2$). In the pivotal mechanism, each agent i reports his net value \tilde{v}_i and the outcome is determined as follows:

Rule 1: if $\tilde{v}_1 + \tilde{v}_2 \geq 0$, then the public good is produced, and if not, then it is not produced; and

Rule 2: each agent i must pay the pivotal tax t_i

$$t_i = \begin{cases} 0 & \text{if (i) } \tilde{v}_j(\tilde{v}_1 + \tilde{v}_2) > 0 \text{ or (ii) } \tilde{v}_j > 0 \text{ and } \tilde{v}_1 + \tilde{v}_2 = 0 \\ |\tilde{v}_j| & \text{otherwise} \end{cases}$$

where $j \neq i$.

That is, an agent pays the amount needed to implement his preferred outcome if his report is pivotal and changes the chosen outcome.

First, let $(v_1, v_2) = (5, -4)$ be the true net value vector. Figure 1-(a) shows that the set of Nash equilibria is approximately a half of the two dimensional area. Notice that the public good should be produced because the sum of the net values of the public good is positive. The upper-right part of the set of Nash equilibria is good since constructing the public good is recommended. However, the lower-left part of the set of Nash equilibria is *bad* since producing the public good is not recommended.

Second, let $(v_1, v_2) = (5, 5)$ be the true net value vector. In this case, both agents want to construct the public good. However, Figure 1-(b) shows the area of bad Nash equilibria is still large. Saijo et al. (2003) generalize this negative result to the case with any arbitrary finite numbers of public projects and agents.

 Figure 1 is around here.

Example 2: The second price auction (Vickrey (1961)).

Consider a two-agent model with an indivisible good. Agent i 's true value of the good is $v_i \geq 0$ if she receives it, and her true value is 0 otherwise ($i = 1, 2$). Let $(\tilde{v}_1, \tilde{v}_2)$ be a reported value vector. The second price auction consists of two rules:

Rule 1: if $\tilde{v}_i > \tilde{v}_j$, then agent i receives the good and pays \tilde{v}_j ($i, j = 1, 2; i \neq j$); and

Rule 2: if $\tilde{v}_1 = \tilde{v}_2$, then agent 1 receives the good and pays \tilde{v}_2 .

Let $(v_1, v_2) = (7, 5)$ be the true value vector. Figure 2 shows that the set of Nash equilibria is quite large. Notice that agent 1 should receive the good because her value is greater than agent 2's. The lower-right part of the set of Nash equilibria is good since agent 1

receives the good. However, the upper-left part of the set of Nash equilibria involving overbidding is bad since agent 2 receives the good.

Figure 2 is around here.

We do not dispute the possibility that, in practice, some confused bidders may fail to recognize their dominant strategy because it is not transparent (e.g., Harstad, 2000). However, our key observation is that the Nash equilibrium areas shown in Figure 2 indicate the possibility of systematic rather than random deviations from the dominant strategy equilibrium.

4. Secure Implementation in Public Good Economies

The previous section presented two examples drawn from many strategy-proof mechanisms that may have “bad” Nash equilibria. They implement the social choice function (SCF) in dominant strategies, but not in Nash equilibria. Saijo et al. (2003) introduce a new concept of implementation, called secure implementation, which does not share this shortcoming.

We introduce notation and definitions here to describe the concept of secure implementation in the context of public good economies with two agents and quasi-linear preferences. Denote the set of feasible allocations by

$$A = \{(y, t_1, t_2) \mid y \in Y, t_1, t_2 \in \mathfrak{R}\},$$

where $Y \subseteq \mathfrak{R}$ is a production possibility set, $y \in Y$ is an output level of a public good, and t_i is a transfer of a private good to agent i . For simplicity, we assume that there is no cost involved in producing y . Each agent i 's utility function, $u_i: A \rightarrow \mathfrak{R}$, is selfish and quasi-linear:

$$u_i(y, t_1, t_2) = u_i(y, t_i) = v_i(y) + t_i, \quad i = 1, 2.$$

The class of valuation functions, $v_i: Y \rightarrow \mathfrak{R}$, admissible for agent i is denoted by V_i . Let $v = (v_1, v_2) \in V \equiv V_1 \times V_2$ be a valuation profile.

A *social choice function* (SCF) is a function $f: V \rightarrow A$ that associates with every list of valuation functions, $v \in V$, a unique feasible allocation $f(v)$ in A . The allocation $f(v)$ is said to be *f-optimal* for v .

A *mechanism* (or *game form*) is a function $g: S_1 \times S_2 \rightarrow A$ that assigns to every $(s_1, s_2) \in S_1 \times S_2$ a unique element of A , where S_i is the *strategy space of agent i*. For a strategy profile $s = (s_1, s_2) \in S_1 \times S_2$, the outcome of g for the profile s is denoted by $g(s) = (y^g(s), t^g(s))$, where $y^g(s)$ is the level of the public good and $t^g(s) = (t_1^g(s), t_2^g(s))$ is the transfer vector.

The strategy profile $s = (s_1, s_2) \in S_1 \times S_2$ is a *Nash equilibrium of g at v* $\in V$ if

$$v_1(y^g(s_1, s_2)) + t_1^g(s_1, s_2) \geq v_1(y^g(s'_1, s_2)) + t_1^g(s'_1, s_2) \text{ for all } s'_1 \in S_1, \text{ and}$$

$$v_2(y^g(s_1, s_2)) + t_2^g(s_1, s_2) \geq v_2(y^g(s_1, s'_2)) + t_2^g(s_1, s'_2) \text{ for all } s'_2 \in S_2.$$

Let $N_A^g(v)$ be the set of Nash equilibrium allocations of g at v , i.e., $N_A^g(v) \equiv \{(y, t_1, t_2) \in A \mid \text{there exists a Nash equilibrium at } v, s \in S, \text{ such that } g(s) = (y, t_1, t_2)\}$.

The strategy profile $s = (s_1, s_2) \in S_1 \times S_2$ is a *dominant strategy equilibrium of g at v* $\in V$ if

$$v_1(y^g(s_1, s'_2)) + t_1^g(s_1, s'_2) \geq v_1(y^g(s'_1, s'_2)) + t_1^g(s'_1, s'_2) \text{ for all } s'_1 \in S_1 \text{ and } s'_2 \in S_2; \text{ and}$$

$$v_2(y^g(s'_1, s_2)) + t_2^g(s'_1, s_2) \geq v_2(y^g(s'_1, s'_2)) + t_2^g(s'_1, s'_2) \text{ for all } s'_1 \in S_1 \text{ and } s'_2 \in S_2.$$

Let $D_A^g(v)$ be the set of dominant strategy equilibrium allocations of g at v , i.e., $D_A^g(v) \equiv \{(y, t_1, t_2) \in A \mid \text{there exists a dominant strategy equilibrium at } v, s \in S, \text{ such that } g(s) = (y, t_1, t_2)\}$.

Definition 1. The mechanism g *implements the SCF f in dominant strategy equilibria* if for all $v \in V$, $f(v) \in D_A^g(v)$. f is *implementable in dominant strategy equilibria* if there exists a mechanism which implements f in dominant strategy equilibria.

Definition 2. The mechanism g *securely implements the SCF f* if for all $v \in V$, $f(v) \in D_A^g(u) = N_A^g(u)$.⁴ The SCF f is *securely implementable* if there exists a mechanism which securely implements f .

⁴ Secure implementation is identical to *double* implementation in dominant strategy equilibria and Nash equilibria. It was Maskin (1979) who first introduced the concept of double implementation. See also Yamato (1993).

Dominant strategy implementation requires that for every possible preference profile, the dominant strategy equilibrium outcome coincides with the f -optimal outcome. In addition to this requirement, secure implementation demands that there be no Nash equilibrium outcome other than the dominant strategy equilibrium outcome.

Saijo et al. (2003) characterize the class of securely implementable SCF's using two conditions. The first condition is strategy-proofness. The allocation recommended by the SCF f for the profile $v = (v_1, v_2)$ is denoted by $f(v) = (y^f(v), t^f(v))$, where $y^f(v)$ is the level of the public good and $t^f(v) = (t_1^f(v), t_2^f(v))$ is the transfer vector.

Definition 3. The SCF f is *strategy-proof* if

$$v_1(y^f(v_1, \tilde{v}_2)) + t_1^f(v_1, \tilde{v}_2) \geq v_1(y^f(\tilde{v}_1, \tilde{v}_2)) + t_1^f(\tilde{v}_1, \tilde{v}_2) \text{ for all } \tilde{v}_1 \in V_1 \text{ and } \tilde{v}_2 \in V_2; \text{ and}$$

$$v_2(y^f(\tilde{v}_1, v_2)) + t_2^f(\tilde{v}_1, v_2) \geq v_2(y^f(\tilde{v}_1, \tilde{v}_2)) + t_2^f(\tilde{v}_1, \tilde{v}_2) \text{ for all } \tilde{v}_1 \in V_1 \text{ and } \tilde{v}_2 \in V_2.$$

By the Revelation Principle (Gibbard, 1973), strategy-proofness is necessary for dominant strategy implementation, and therefore also for secure implementation. However, the following additional condition, called the *rectangular property*, is necessary for secure implementation.

Definition 4. The SCF f satisfies the *rectangular property* if for all $v, \tilde{v} \in V$, if

$$v_1(y^f(v_1, \tilde{v}_2)) + t_1^f(v_1, \tilde{v}_2) = v_1(y^f(\tilde{v}_1, \tilde{v}_2)) + t_1^f(\tilde{v}_1, \tilde{v}_2) \text{ and}$$

$$v_2(y^f(\tilde{v}_1, v_2)) + t_2^f(\tilde{v}_1, v_2) = v_2(y^f(\tilde{v}_1, \tilde{v}_2)) + t_2^f(\tilde{v}_1, \tilde{v}_2),$$

then $f(v_1, v_2) = f(\tilde{v}_1, \tilde{v}_2)$.

Saijo et al. (2003) show that the rectangular property is necessary and sufficient for sure implementation:⁵

Theorem 1. *An SCF is securely implementable if and only if it satisfies strategy-proofness and the rectangular property.*

By applying Theorem 1, Saijo et al. (2003) find that no strategy-proof and efficient SCF is securely implementable if public goods are discrete and no restrictions are placed on the set of admissible valuation functions.⁶

Consider an SCF f satisfying the efficiency condition on the public good provision:

$$(4.1) \quad y^f(v_1, v_2) \in \arg \max_{y \in Y} [v_1(y) + v_2(y)] \text{ for all } (v_1, v_2) \in V.$$

The following result is well known:

⁵ To see why the rectangular property is necessary for secure implementation intuitively, suppose that the direct revelation mechanism $g = f$ securely implements the SCF f . Let $n = 2$ and (v_1, v_2) be the true preference profile.

Suppose $u_1(f(v_1, \tilde{v}_2)) = u_1(f(\tilde{v}_1, \tilde{v}_2))$, i.e.,

$$(*) \quad v_1(y^f(v_1, \tilde{v}_2)) + t_1^f(v_1, \tilde{v}_2) = v_1(y^f(\tilde{v}_1, \tilde{v}_2)) + t_1^f(\tilde{v}_1, \tilde{v}_2).$$

In other words, agent 1 is indifferent between reporting the true preference v_1 and reporting another preference \tilde{v}_1 when agent 2's report is \tilde{v}_2 . Since reporting v_1 is a dominant strategy by strategy-proofness, it follows from (*) that

$$v_1(y^f(\tilde{v}_1, \tilde{v}_2)) + t_1^f(\tilde{v}_1, \tilde{v}_2) = v_1(y^f(v_1, \tilde{v}_2)) + t_1^f(v_1, \tilde{v}_2) \geq v_1(y^f(v'_1, \tilde{v}_2)) + t_1^f(v'_1, \tilde{v}_2) \text{ for all } v'_1 \in V_1,$$

that is, reporting \tilde{v}_1 is one of agent 1's best responses when agent 2 reports \tilde{v}_2 .

Next suppose that $u_2(f(\tilde{v}_1, v_2)) = u_2(f(\tilde{v}_1, \tilde{v}_2))$, i.e.,

$$(**) \quad v_2(y^f(\tilde{v}_1, v_2)) + t_1^f(\tilde{v}_1, v_2) = v_2(y^f(\tilde{v}_1, \tilde{v}_2)) + t_2^f(\tilde{v}_1, \tilde{v}_2).$$

By using an argument similar to the above, it is easy to see that $v_2(y^f(\tilde{v}_1, \tilde{v}_2)) + t_2^f(\tilde{v}_1, \tilde{v}_2) = v_2(y^f(\tilde{v}_1, v_2)) + t_2^f(\tilde{v}_1, v_2) \geq v_2(y^f(\tilde{v}_1, v'_2)) + t_2^f(\tilde{v}_1, v'_2)$ for all $v'_2 \in V_2$,

that is, reporting \tilde{v}_2 is one of agent 2's best responses when agent 1 reports \tilde{v}_1 . Therefore, $f(\tilde{v}_1, \tilde{v}_2) =$

$(y^f(\tilde{v}_1, \tilde{v}_2), t^f(\tilde{v}_1, \tilde{v}_2))$ is the Nash equilibrium outcome. Moreover, $f(v_1, v_2) = (y^f(v_1, v_2), t^f(v_1, v_2))$ is the dominant strategy outcome, and by secure implementability, the dominant strategy outcome coincides with the Nash equilibrium outcome. Accordingly we conclude that $f(v_1, v_2) = f(\tilde{v}_1, \tilde{v}_2)$ if (*) and (**) holds.

⁶ We say that an SCF is efficient if it produces the efficient public goods level, i.e., if (4.1) is satisfied. We do not require "budget balance" in the sense that the tax revenue equals the cost of producing the public good. As is well known, (4.1) and strategy proofness will in general force the budget to be unbalanced.

Proposition 1 (Clarke (1971), Groves (1973), Green and Laffont (1979)). An SCF f satisfying (4.1) is implementable in dominant strategy equilibria if and only if f satisfies

$$(4.2) \quad t_1^f(v_1, v_2) = v_2(y^f(v_1, v_2)) + h_1(v_2), \quad t_2^f(v_1, v_2) = v_1(y^f(v_1, v_2)) + h_2(v_1) \quad \forall (v_1, v_2) \in V,$$

where h_i is some arbitrary function which does not depend on v_i .

A direct revelation mechanism satisfying (4.1) and (4.2) is called a *Groves-Clarke mechanism*. Proposition 1 says that we can focus on the class of Groves-Clarke mechanisms for implementation of an efficient SCF in dominant strategy equilibria. However, Saijo et al. (2003) show that if we do not put any restrictions on V , and if Y is a finite set, then for any mechanism implementing an efficient SCF in dominant strategy equilibria, the set of Nash equilibrium outcomes is strictly larger than the set of dominant strategy equilibrium outcomes.

Results are different if V contains only single-peaked preferences and y is a continuous variable. In this case strategy-proof and efficient SCF's are securely implementable by Groves-Clarke mechanisms. Suppose that $Y = \mathfrak{R}$ and for $i = 1, 2$,

$$V_i = \{v_i: \mathfrak{R} \rightarrow \mathfrak{R} \mid v_i(y) = -(y - r_i)^2, r_i \in \mathfrak{R}\},$$

where r_i is agent i 's most preferred level of the public good. We can represent these single-peaked preferences by the r_i instead of the v_i . The optimal output level of the public good satisfying (4.1) is given by $y(r_1, r_2) = (r_1 + r_2)/2$.

For this case any SCF f meeting (4.1) and (4.2) satisfies the rectangular property and is therefore securely implementable (Saijo et al., 2003).

Consider an example that will be used in our experimental design later, in which $h_i = 0$. Then,

$$\begin{aligned} u_1(\tilde{r}_1, \tilde{r}_2) &= v_1(y(\tilde{r}_1, \tilde{r}_2)) + t_1(\tilde{r}_1, \tilde{r}_2) = -((\tilde{r}_1 + \tilde{r}_2)/2 - r_1)^2 - ((\tilde{r}_1 + \tilde{r}_2)/2 - \tilde{r}_2)^2 \\ &= -\{(\tilde{r}_1 - r_1)^2 + (\tilde{r}_2 - r_1)^2\} / 2 \end{aligned}$$

where r_1 is player 1's true peak and $(\tilde{r}_1, \tilde{r}_2)$ is a vector of reported peaks. Clearly agent 1's payoff is maximized at r_1 . Since the payoff function is quadratic, no other maximizers exist. Furthermore, the payoff is maximized at r_1 regardless \tilde{r}_2 . Figure 3 shows agent 1's payoff when $r_1 = 12$. If $\tilde{r}_2 = 4$, the maximizer is a , and if $\tilde{r}_2 = 12$, it is b . Both are maximized at $r_1 = 12$. Therefore, the best response curve is a line parallel to the \tilde{r}_2 axis. This indicates that truth-telling is *the* dominant strategy. In fact, it is *strictly* dominant. However, this is true only as long as the public goods level is continuously variable. In our experiment, we will discretize the public goods level and the payoff functions, and truth-telling will not be strictly dominant even though preferences are single-peaked.⁷ However, with single-peaked preferences implementation will still be *secure*, because there will be a unique dominant strategy equilibrium which is also a unique Nash equilibrium (Treatment S). When preferences are not single peaked, there will be multiple Nash equilibria and implementation is not secure (Treatment P).

Figure 3 is around here.

5. The Experiment

Our experiment studies the pivotal mechanism and a Groves-Clarke mechanism with single-peaked preferences. It consisted of four sessions with 20 subjects each (80 total subjects). We conducted two sessions in Treatment P that corresponded to the pivotal mechanism and two sessions in Treatment S that corresponded to a Groves mechanism with single-peaked preferences.

5.1 Design

We conducted two sessions (one P and one S) at Tokyo Metropolitan University during June of 1998 and two sessions (one P and one S) at Purdue University during February of 2003.

⁷ In general, with a discrete public good, single-peaked preferences will not assure the existence of a strictly dominant strategy. However, secure implementation will be assured.

Each session took approximately one hour to complete.

Treatment P implements the pivotal mechanism for a two-person group. The net true value vector (v_1, v_2) is equal to $(-6, 8)$ if a binary public good is produced and $(v_1, v_2) = (0, 0)$ otherwise. The public good should be produced since $v_1 + v_2 \geq 0$. Let the strategy space of type 1 be the set of integers from -22 to 2, and the strategy space of type 2 be the set of integers from -4 to 20. According to the rules of the pivotal mechanism described in Section 3, we can construct the payoff matrix of types 1 and 2 as Tables C-1 and C-2 shown in Appendix C.

The payoff tables that we actually distributed to subjects in Treatment P were Tables 1 and 2 whose basic structures were the same as Tables C-1 and C-2. But we modified Tables C-1 and C-2 as follows. First, we changed the names of strategies. Type 1's strategy "-22" was renamed "1", "-21" was renamed "2", and so on. Similarly, type 2's strategy "-4" was renamed "1", "-3" was renamed "2", and so on. Second, we employed a linear transformation of the valuation functions: $14v_1 + 294$ for type 1 and $14v_2 + 182$ for type 2. Of course, the equilibrium regions shown on these versions of the tables were not displayed to subjects.

 Tables 1 and 2 are around here.

Treatment S is the same as Treatment P except for the payoff tables. The payoff tables for Treatment S are based on the following model of a Groves mechanism with single-peaked preferences with two players. Suppose that the true valuation functions of agent types 1 and 2 are respectively $v_1(y) = -(y - 12)^2$ and $v_2(y) = -(y - 17)^2$, where $y \in \mathfrak{R}_+$ is the level of a public good. Each type reports his most preferred level of the public good called a peak. Given a vector of reported peaks $(\tilde{r}_1, \tilde{r}_2)$, the level of the public good, $y(\tilde{r}_1, \tilde{r}_2)$, and the transfer to type i , $t_i(\tilde{r}_1, \tilde{r}_2)$, are determined by a Groves mechanism: $y(\tilde{r}_1, \tilde{r}_2) = (\tilde{r}_1 + \tilde{r}_2) / 2$ and $t_i(\tilde{r}_1, \tilde{r}_2) = -((\tilde{r}_1 + \tilde{r}_2) / 2 - \tilde{r}_j)^2$, $i, j = 1, 2; j \neq i$. The payoff functions are therefore given by

$$\begin{aligned} v_1(y(\tilde{r}_1, \tilde{r}_2)) + t_1(\tilde{r}_1, \tilde{r}_2) &= -((\tilde{r}_1 + \tilde{r}_2) / 2 - 12)^2 - ((\tilde{r}_1 + \tilde{r}_2) / 2 - \tilde{r}_2)^2, \\ v_2(y(\tilde{r}_1, \tilde{r}_2)) + t_2(\tilde{r}_1, \tilde{r}_2) &= -((\tilde{r}_1 + \tilde{r}_2) / 2 - 17)^2 - ((\tilde{r}_1 + \tilde{r}_2) / 2 - \tilde{r}_1)^2. \end{aligned}$$

Let the strategy space of each type be the set of integers from 0 to 24. The payoff table for types 1 and 2 are is given by Tables C-3 and C-4 in Appendix C.

The payoff tables used in Treatment S were Tables 3 and 4 whose basic structures were the same as Tables C-3 and C-4, modified as follows. First, we changed the names of strategies: strategy "0" was renamed "1", "1" was renamed "2", and so on. Second, we employed a linear transformation of the payoff functions: $10v_i / 14 + 218.5$ for $i = 1, 2$.

Tables 3 and 4 are around here.

Note that because we discretized the possible levels in the payoff tables and rounded payoffs to the nearest whole number, neither player type has a strictly dominant strategy. Therefore, Treatments S and P cannot be differentiated in terms of strictly dominant strategies. However, only Treatment S involves a secure mechanism.

5.2 Procedures

The sessions in Japan and in the United States involved a variety of procedural differences. They were not intended to replicate the same experimental conditions, but instead were useful to evaluate the robustness of our findings to different subject pools and procedures. Most notably, the sessions in Japan were run "by hand" with pen and paper, and the sessions in the U.S. were computerized using zTree (Fischbacher, 1999). If we had observed significant differences across experiment sites, then we would not be able to identify the source of those differences without further experimentation. Fortunately, as we show in the Appendix B the data do not indicate any meaningful statistically significant differences across sites within either mechanism treatment.

In the Japan sessions the twenty subjects were seated at desks in a relatively large room and had identification numbers assigned randomly. These ID numbers were not publicly displayed, however, so subjects could not determine who had which number. In the U.S. sessions the twenty subjects were seated at computer stations in the Vernon Smith Experimental

Economics Laboratory that were separated with visual partitions. In every period, each of the type 1 subjects was paired with one of the type 2 subjects. The pairings were determined in advance by experimenters so as not to pair the same two subjects more than once (“strangers”). Each subject received written instructions, a record sheet, a payoff table, and (in the Japan sessions only) information transmission sheets. Instructions (see Appendix A) were also given by tape recorder in Japan and were read aloud by the experimenter in the U.S. Each subject chose her number from an integer between 1 and 25 by looking at her own payoff table only.⁸ No subject knew the payoff table of the other type. Moreover, we provided no explanation regarding the rules of the mechanisms or how the payoff tables were constructed.

After deciding which number she chose, each subject marked the number on an information transmission sheet (Japan) or typed in her number on her computer (U.S.). Experimenters collected these information transmission sheets and then redistributed them to the paired subjects in Japan. The computer network handled the message transmission in the U.S. Each period, subjects in both countries were asked to fill out the reasons why they chose these numbers. After learning the paired subject’s choice, subjects calculated their payoffs from the payoff tables (Japan) or verified the computer-calculated payoffs (U.S.). Record sheets were identical (except for the language translation, of course) at the two sites. These steps were repeated for eight periods in Japan and for ten periods in the U.S. Recall that subjects were never paired together for more than one period.

In the Japan sessions the mean payoff per subject was 1677 yen in Treatment S and it was 1669 yen in Treatment P. In the U.S. sessions the mean payoff per subject was \$21.04 in Treatment S and it was \$20.35 in Treatment P.

⁸ We required subjects to examine their payoff table for ten minutes before we began the real periods.

6. Results

6.1 Treatment P

Tables 1 and 2 specify the dominant strategy equilibria and the other Nash equilibria for Treatment P. Type 1's dominant strategies are 16 and 17, and type 2's dominant strategies are 12 and 13. Notice that the dominant strategy equilibria (16, 12) and (16, 13) are Pareto-dominated by the dominant strategy equilibria (17, 12) and (17, 13). In Table 1, the lower-right region of Nash equilibria is good since the public good is produced. The upper-left region of Nash equilibria is bad since the public good is not produced. The number of good Nash equilibria is 162, while the number of bad Nash equilibria is 165. Implementation is clearly not secure.

Since each period had 20 pairs of players and each session had 8 or 10 periods, we have 180 pairs of data. Denote each pair by (x_1, x_2) where x_i is a number chosen by a subject of type i , $i = 1, 2$. Figure 4 shows the frequency distribution of all data in Treatment P. The maximum frequency pair was (16,12) with 34 pairs, the second was (16, 13) with 27 pairs, the third was (17,13) with 19 pairs, and the fourth was (17, 12) with 10 pairs. The total frequency of the four dominant strategy equilibria (16,12), (16,13), (17,12), and (17,13) was 90—exactly one-half of the outcomes. The frequency of Pareto-dominated dominant strategy equilibria (16,12) and (16,13) was 61, while the frequency of the dominant strategy equilibria (17,12) and (17,13) was 29.⁹ Fifty-nine other outcomes were good Nash equilibria other than dominant strategy equilibria. Only one pair in one period played a bad Nash equilibrium. Although nearly half (298/621) of the strategy pairs shown in Tables 1 and 2 that are not dominant strategy equilibrium outcomes are not Nash equilibria, only one-third (30/90) of the observed non-dominant-strategy outcomes were not Nash equilibria. This suggests that deviations from the

⁹ We are puzzled by the greater frequency of Pareto dominated dominant strategy equilibria. Seventy-two percent of the dominant strategies played by Type 1 players were 16 rather than 17. The greater frequency of 16 declines in later periods, however, and only in periods 1 and 3 is 16 significantly more frequent than 17 at the 5-percent level (two-tailed) according to a binomial test.

dominant strategy equilibria are not random, but are instead more likely to correspond to Nash equilibria.

Figure 4 is around here

We conducted period by period tests of the hypothesis that the mean choice is equal to a dominant strategy (16 or 17 for type 1 and 12 or 13 for type 2). A nonparametric Wilcoxon signed rank test rejects the hypothesis that type 1 subjects play the dominant strategy of 17 in five out of ten periods (periods 1, 2, 3, 7 and 8), but this test never rejects the null hypothesis that subjects on average play the dominant strategy of 16 (two-tailed test, five-percent significance level). Similarly, this nonparametric test rejects the hypothesis that type 2 subjects play the dominant strategy of 12 in eight out of ten periods (periods 1, 2, 3, 5, 6, 7, 8 and 9), but this test never rejects the null hypothesis that subjects on average play the dominant strategy of 13.

These Treatment P results lead to the following observations:

Observation 1:

- (a) *The frequency of dominant strategy equilibria was 50% across all periods in Treatment P.*
- (b) *Subjects played Pareto-dominated dominant strategy equilibria about twice as frequently as Pareto-superior dominant strategy equilibria in Treatment P.*
- (c) *The data do not reject the hypothesis that subjects play a dominant strategy on average for either type in any period in Treatment P.*
- (d) *Almost all (98%) of the observed Nash equilibria other than the dominant strategy equilibria were good Nash equilibria that recommended funding of the public good.*

6.2 Treatment S

A unique dominant strategy equilibrium exists in Tables 3 and 4: 13 for type 1 and 18 for type 2. Notice that there are no other Nash equilibria. Thus, even though the dominant

strategy is not strictly dominant, the implementation is secure.

Figure 5 shows the frequency distribution of all data in Treatment S. The maximum frequency pair was the dominant strategy equilibrium (13,18) with 146 of the 180 outcomes. Pairs played no other single outcome more than 4 times.

Figure 5 is around here

We conducted period by period tests of the hypothesis that the mean number equals the dominant strategy (13 for type 1 and 18 for type 2). A Wilcoxon signed rank test never rejects the dominant strategy equilibrium hypothesis for any type in any period.

Summarizing the above results, we have the following:

Observation 2:

- (a) *The frequency of dominant strategy equilibrium was 81% across all periods in Treatment S.*
- (b) *The data do not reject the hypothesis that subjects choose the dominant strategy on average for either type in any period in Treatment S.*

6.3 Comparing the Two Mechanisms

Here we compare the frequency that subjects play dominant strategies and that pairs implement dominant strategy equilibria in the two mechanisms. Recall that an advantage of our experimental design is that we can compare these two mechanisms while holding constant their complexity. We did not present to subjects any explanation on the rules of a mechanism, and instead we simply used payoff tables to explain the relationship between choices and outcomes.

Figure 6 displays the rates that subjects play dominant strategies separately for all periods. Individuals are more likely to play dominant strategies in Treatment S than in Treatment P according to Fisher's exact test in 7 out of 10 periods (periods 2, 6, 7, 8 and 9 at the 5% significance level, and periods 4 and 5 at the 10% significance level). A more powerful

parametric test is possible by pooling the data across periods. Since individual subjects contribute an observation for each period it is appropriate to model the panel nature of the data. We do this with a subject random effect specification for the error term.¹⁰ Column 1 of Table 5 reports a probit model of the likelihood that the subject selects a dominant strategy. The positive and significant dummy variable for the mechanism treatment indicates that subjects are more likely to play a dominant strategy in the secure mechanism.¹¹

Figure 6 and Table 5 are around here

Figure 7 shows that the differences in the individual dominant strategy rates are magnified for the pair rates. Pairs are more likely to play a dominant strategy equilibrium in Treatment S according to Fisher's exact test in 8 out of 10 periods (periods 2, 4, 6, 7, 8 and 9 at the 5% significance level, and periods 3 and 5 at the 10% significance level). Column 2 of Table 5 reports a probit model of the likelihood that pairs play a dominant strategy equilibrium, pooling across periods. A random subject effect specification is not possible since the composition of the individuals in each pair changes each period. But we include a dummy variable for the Purdue sessions to capture any (fixed effect) differences across sessions, and this variable is not significantly different from zero. The mechanism treatment dummy variable is highly significant, however, indicating the substantially greater frequency of dominant strategy equilibrium play in Treatment S. Recall that neither Treatment S nor Treatment P have strictly dominant strategies, but only Treatment S involve a secure mechanism.

¹⁰ Session rather than subject effects provide similar results, and even more statistically significant estimated mechanism treatment effects.

¹¹ Recall that subjects also indicated the reasons for their choices on their record sheets and in a post-experiment questionnaire. We reviewed their responses and found that more individual subjects provided explanations that were clearly identifiable as dominant strategy arguments (e.g., "This is the highest payoff column no matter what the other person chooses.") in Treatment S (23 individuals) than in Treatment P (13 individuals). This difference is statistically significant according to Fisher's exact test (p -value=0.021).

Figure 7 also illustrates the frequency of Nash equilibrium play at each period for Treatment P, which is much higher than that of dominant strategy equilibrium play. Because almost all Nash equilibria were “good”, this suggests that the pivotal mechanism was not for dominant strategy implementation, but for Nash implementation in our experiment, although the mechanism may produce “bad” Nash equilibria in theory.

Figure 7 is around here

Summarizing the above results, we have the following:

Observation 3:

- (a) *Individuals play dominant strategies significantly more frequently in Treatment S than in Treatment P.*
- (b) *Pairs implement dominant strategy equilibria significantly more frequently in Treatment S than in Treatment P.*

7. Conclusion

Recent experimental and theoretical findings have raised serious questions about the viability of dominant strategy mechanisms. A possible solution is the notion of *secure implementation* introduced in Saijo, Sjöström and Yamato (2003). Motivated by this theoretical concept this paper presents an experimental study of the pivotal mechanism and a Groves-Clarke mechanism with single-peaked preferences. Both mechanisms are strategy-proof since they implement truth-telling as dominant strategy equilibria. But the pivotal mechanism has other Nash equilibria that differ from the dominant strategy equilibria. Although the design simplifies both mechanisms equally with payoff tables, players adopted dominant strategies significantly less often in the pivotal mechanism.

This experiment illustrates how a mechanism that is not secure may not yield the desired outcome. Indeed, in the non-secure mechanism the players failed to use their dominant strategies about half of the time. This did not have negative welfare implications, because the public good was anyway funded 98% of the time. We believe that this outcome was fortuitous, and that there is no reason in general why players who do not play their dominant strategies will end up funding the public good in the right amount in a non-secure mechanism. In contrast, we are optimistic about the performance of secure mechanisms, because they are more likely to induce truth-telling.

We believe that this point is relevant for practical mechanism design. An obvious application is the second price (Vickrey) auction. For certain information conditions – most clearly in a complete information setting – other “bad” Nash equilibria exist in this auction that do not correspond to the (efficient) dominant strategy, truth-telling equilibrium. Most proponents of this auction institution have not acknowledged this shortcoming. Before making predictions regarding how this institution might perform in the field, it would be valuable to conduct laboratory experiments with the information conditions that admit these other inefficient Nash equilibria. We suspect that the second price auction and many other strategy-proof mechanisms may not function as elegantly as designed on the theorist’s blackboard, due in part to the frequent existence of non-dominant strategy Nash equilibria.

In practice, mechanisms cannot be too complex, due to the finite information processing capacity of the players. It turns out that requiring secure implementation does not lead to more complex mechanisms: attention can be restricted to revelation mechanisms without loss of generality (Saijo, Sjöström and Yamato (2003)). Unfortunately, this makes it quite difficult to test the hypothesis that secure mechanisms perform better than non-secure ones, at least in the public goods economy. Indeed, by Proposition 1, efficiency and strategy-proofness essentially pin down the revelation mechanism. Consequently, we cannot compare the performance of two efficient strategy-proof mechanisms, one that is secure and one that is not, in a fixed environment. The environment (specifically, the set of valuation functions) has to vary across

treatments (as in Treatment P versus Treatment S). In our experiment, we do not think this matters too much, because the presentation of the payoff tables was similar in the two treatments. Still, in some economic models it may be possible to make interesting comparisons of the performance of secure versus non-secure mechanisms *in the same environment*. This is left for future experiments.

Appendix A: Experiment Instructions

Instructions

This is an experiment in the economics of strategic decision making. Various agencies have provided funds for this research. If you follow the instructions and make appropriate decisions, you can earn an appreciable amount of money. At the end of today's session, you will be paid in private and in cash.

Overview

In this experiment you will choose a number in each period. You will be paired with one other person each period. The other person will also choose a number simultaneously. Your payoff is determined by the number you choose as well as the number the other person chooses. You can see how much your payoff is by looking at your payoff table. However, your payoff table is different from the payoff table the other person has. An experimenter will choose the person you are paired with from the other participants at random, and the person you are paired with will change each period. You will never be paired with the same person more than once.

This experiment consists of 10 periods. Your earnings are the sum of your payoffs over all 10 periods.

In this experiment, please remember that you cannot talk to anyone but the experimenter. If there is any other talking, the experiment will be stopped at that point. If you have any questions, please ask an experimenter.

First, please confirm the following items on your desk.

- Instructions (this set of papers)
- Pencil or pen
- Payoff Table for Practice
- Record Sheet for Practice

Practice

Please look at the "Payoff Table for Practice." Your payoff table is different from the payoff table of the other people you will be paired with in the actual experiment. However, everyone has the same payoff table in these practice examples.

You will choose one integer number from 1 to 25. Suppose that the number you choose is "18". At the time you choose your number, you do not know which number the other person chooses.

Please write the number “18” into the column of “The Number You Chose” in the “Record Sheet for Practice.”

In the actual experiment you will also type the number you chose on your computer. The computer will transmit your choice to the person you are paired with for the period. In the practice examples, we’ll skip this part. Moreover, in the actual experiment, please fill out why you chose that number into the column of the “Reason for Your Decision” in the record sheet. In practice, we skip this part, too.

Suppose that the person you are paired with also chose “18” in this practice example. Please look at the payoff table. Your payoff is “150” when the number you chose is “18” and the number the other person chose is “18”. Your earnings in this period are equal to the value of your payoff, that is, 150 cents. Please write “18” into the column of “The Number the Other Person Chose” and “150” into the column of “Your Payoff” in the “Record Sheet for Practice” in the Period 1 row.

Payoff Table for Practice

		The number you choose																								
		1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25
The number the other person chooses	1	65	65	66	68	70	72	75	78	81	83	85	87	87	88	87	86	84	82	79	76	73	70	67	66	65
	2	65	66	68	70	73	76	80	83	86	88	90	91	92	92	91	89	87	84	81	77	74	70	68	66	65
	3	66	68	71	74	78	81	85	88	91	94	96	97	97	97	95	93	90	87	83	79	75	71	68	66	65
	4	69	71	75	79	83	87	91	95	98	100	102	103	103	102	100	97	94	90	85	80	76	72	68	66	65
	5	72	76	80	85	90	94	99	102	105	108	109	110	109	107	105	101	97	93	87	82	77	72	68	66	65
	6	77	82	87	92	98	103	107	111	114	116	117	117	116	114	110	106	101	96	90	84	78	73	69	66	65
	7	83	89	95	101	107	112	117	120	123	125	125	125	123	120	116	111	105	99	92	86	79	74	69	66	65
	8	91	98	105	111	117	123	127	131	133	135	135	133	131	127	122	116	110	103	95	88	81	74	69	66	65
	9	101	108	116	123	129	135	139	142	144	145	145	143	139	135	129	122	114	106	98	90	82	75	70	66	65
	10	112	121	128	136	142	148	152	155	157	157	155	153	148	143	136	128	119	110	101	92	84	76	70	66	65
	11	126	135	143	150	157	162	166	169	170	169	167	163	158	151	143	134	125	115	104	94	85	77	71	66	65
	12	141	150	159	167	173	178	182	184	184	182	179	174	168	160	151	141	130	119	108	97	87	78	71	67	65
	13	158	168	177	184	191	195	198	200	199	197	192	186	179	170	159	148	136	124	111	99	88	79	71	67	65
	14	177	187	196	204	210	214	216	217	215	212	206	199	190	180	168	156	142	128	115	102	90	80	72	67	65
	15	199	209	218	225	231	234	236	235	233	228	221	213	202	191	177	163	149	134	119	105	92	81	72	67	65
	16	222	233	241	248	253	256	257	255	251	245	237	227	215	202	187	171	155	139	123	108	94	82	73	67	65
	17	248	258	267	273	277	279	279	276	271	263	254	242	228	213	197	180	162	144	127	111	96	83	73	67	65
	18	276	286	294	299	303	304	302	298	291	282	271	258	242	226	208	189	169	150	131	114	98	84	74	67	65
	19	306	316	323	328	330	330	327	321	313	302	289	274	257	239	219	198	177	156	136	117	100	86	75	67	65
	20	339	348	354	358	359	358	353	346	336	324	309	292	273	252	230	208	185	162	141	120	102	87	75	68	65
	21	374	382	388	390	390	387	381	372	360	346	329	310	289	266	242	218	193	169	146	124	104	88	76	68	65
	22	411	418	423	425	423	418	410	399	386	369	350	329	305	281	255	228	202	176	151	127	107	90	76	68	65
	23	451	457	461	461	457	451	441	428	412	393	372	348	323	296	268	239	210	182	156	131	109	91	77	68	65
	24	493	499	500	499	494	485	473	458	440	419	395	369	341	311	281	250	220	190	161	135	112	92	78	68	65
	25	538	542	542	539	532	521	507	490	469	445	419	390	360	328	295	262	229	197	167	139	114	94	78	68	65

Let us take another example for the second practice period. Suppose you again chose 18. Write this in “The Number You Chose” column for Period 2. But for this example suppose that the person you are paired with chose “5” rather than “18”. Then your payoff is 93 cents in this period. Please write “5” into the column of “The Number the Other Person Chose” and “93” into the column of “Your Payoff.”

Finally, let us go on to a third practice period example. Suppose that you chose "6". Write this in Period 3 on your record sheet. Suppose that the person you are paired with chose "19". Then your payoff is 330 cents in this period. Please write "19" into the column of "The Number the Other Person Chose" and "330" into the column of "Your Payoff."

The sum of your payoffs for these three practice periods is $150+93+330=573$. Please write "573" into the column of "the Sum of Your Payoffs." This is only practice so you will not be paid this amount.

If you have any questions, please raise your hand.

The Actual Experiment

First, please pass the "Payoff Table for Practice" and the "Record Sheet of Practice" back to the experimenter now. Next, we will distribute your payoff table and record sheet for the actual experiment.

Your payoff table is different from the payoff table of the people you will be paired with. You will have 10 minutes to look at the payoff table to understand it before we begin the experiment.

Are there any questions?

Appendix B: Comparison of the Japanese and American Sessions Results

B.1 Treatment P

Average choices, Type 1: The data do not reject the null hypothesis that the average choices are equal in the two countries in any of the 8 individual periods (t -test) at 10% significance level. The data reject this same null hypothesis in two periods at the 10% level (only), periods 4 and 5, using a nonparametric Wilcoxon test. When pooling the data across periods and comparing the choices across countries using a panel data regression (with individual subjects as the random effect for the error term), the data do not reject the null hypothesis that choices are equal in the two countries ($t=1.312$).

Average choices, Type 2: The data do not reject the null hypothesis that the average choices are equal in the two countries in any of the 8 individual periods (t -test) at 10% significance level. The data reject this same null hypothesis in one of the 8 periods (period 1) at the 10% level (only) using a nonparametric Wilcoxon test. When pooling the data across periods and comparing the choices across countries using a panel data regression (with individual subjects as the random effect for the error term), the data do not reject the null hypothesis that choices are equal in the two countries ($t=0.609$).

Rate playing dominant strategy, pooling over types: According to Fisher's exact test, the data reject the null hypothesis that the rate subjects play a dominant strategy is equal in the two countries in one of the 8 individual periods (period 1), at the 5% significance level. Pooling the data across periods using a panel data regression (with individual subjects as the random effect for the error term) in a probit model of the likelihood that subjects play a dominant strategy, we *do* reject the null hypothesis that there is no difference across countries ($t=2.001$). But this is due to the significant difference in the first period only; estimating this same model after dropping the first period, we do not reject the null hypothesis of no differences across countries ($t=1.559$).

B.2 Treatment S

Average choices, Type 1: The data reject the null hypothesis that the average choices are equal in the two countries in only one (period 1) of the 8 individual periods (t -test) at the 10% significance level. The data reject this same null hypothesis in only the same one period (period 1) at the 10% level using a nonparametric Wilcoxon test. When pooling the data across periods and comparing the choices across countries using a panel data regression (with individual subjects as the random effect for the error term), the data do not reject the null hypothesis that choices are equal in the two countries ($t=1.774$).

Average choices, Type 2: The data do not reject the null hypothesis that the average choices are equal in the two countries in any of the 8 individual periods at 10% significance level, using either a t -test or a nonparametric Wilcoxon test. When pooling the data across periods and comparing the choices across countries using a panel data regression (with individual subjects as the random effect for the error term), the data do not reject the null hypothesis that choices are equal in the two countries ($t=1.295$).

Rate playing dominant strategy, pooling over types: According to Fisher's exact test, the data do not reject the null hypothesis that the rate subjects play a dominant strategy is equal in the two countries in any of the 8 individual periods at the 10% significance level. Pooling the data across periods using a panel data regression (with individual subjects as the random effect for the error term) in a probit model of the likelihood that subjects play a dominant strategy, we do not reject the null hypothesis of no differences across countries ($t=0.029$).

B.3 Summary

Of the 32 (=2 types * 2 treatments * 8 periods) period-by-period tests comparing individual choices across sites, we reject the null of no country difference in 1 out of 32 (t -test) and 3 out of 32 (Wilcoxon test) at the 10% significance level. We expect about 10% rejections

(about 3 out of 32) at the 10% significance level if the null is true, exactly as we observe. This leads us to the conclusion that there are no country differences. The panel data regressions pooling across periods confirm this conclusion.

Of the 16 (=2 treatments * 8 periods) tests that pool across types, we reject the null that subjects play the dominant strategy at different rates across countries in 1 out of the 16 cases. The panel data regressions confirm that the only difference across sites is in period 1 of treatment P.

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	(1) Individuals play dominant strategies	(2) Pairs play dominant strategy equilibrium
Dummy variable=1 for Treatment S	0.720** (0.346)	0.887** (0.143)
Dummy variable=1 for sessions at Purdue		0.170 (0.142)
Intercept	1.236** (0.266)	-0.095 (.0122)
$\rho = \sigma_u^2 / (\sigma_v^2 + \sigma_u^2)$ (random effects significance)	0.627** (0.069)	
Observations	720	360
Log-likelihood	-247.2	-211.3
Restricted log-likelihood	-344.5	-231.8

Notes: Standard errors shown in parentheses. ** denotes significantly different from zero at five-percent. Model in column (1) is estimated with a random subjects effect error term $u_i + v_{it}$.

Table 5. Probit Models of Individual and Pair Dominant Strategy Play

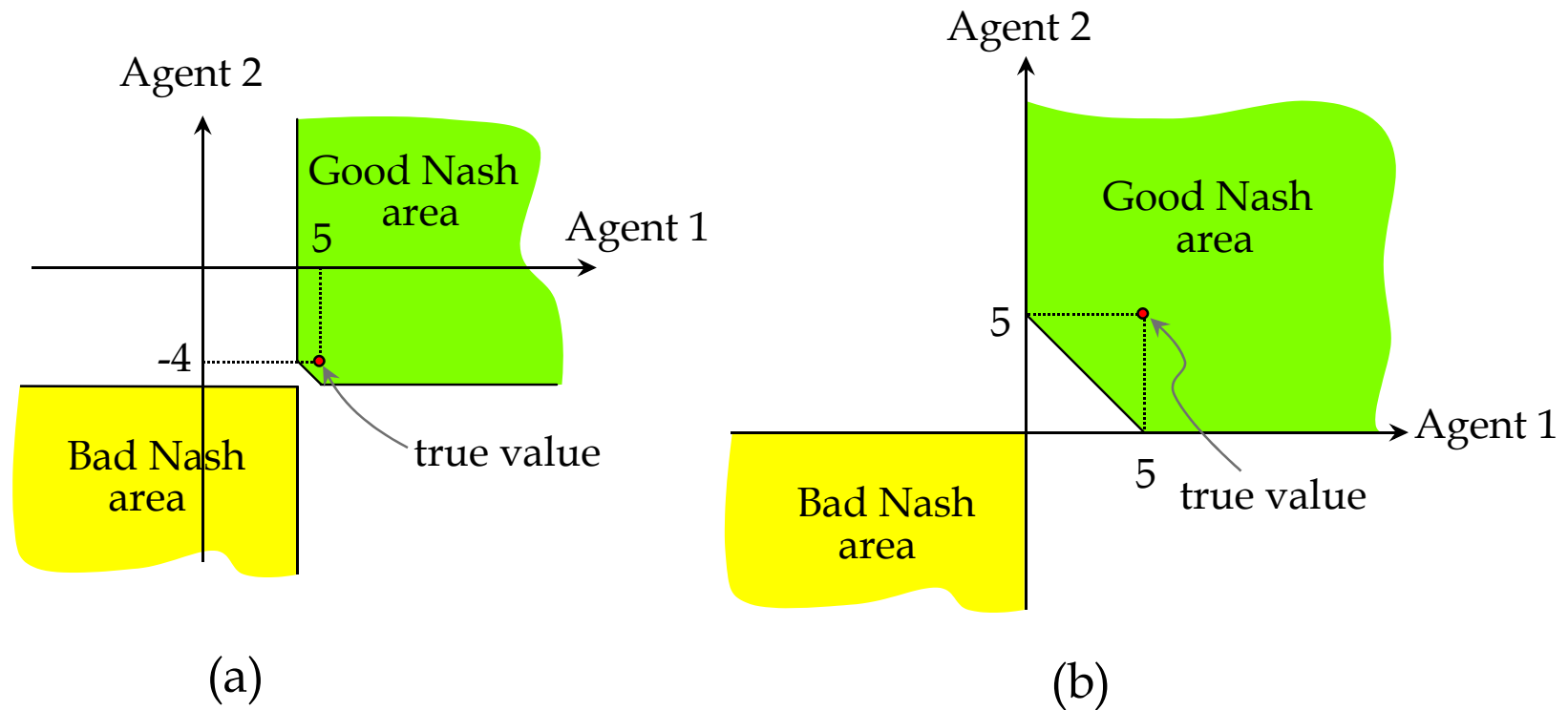


Figure 1: Equilibria of the Pivotal Mechanism

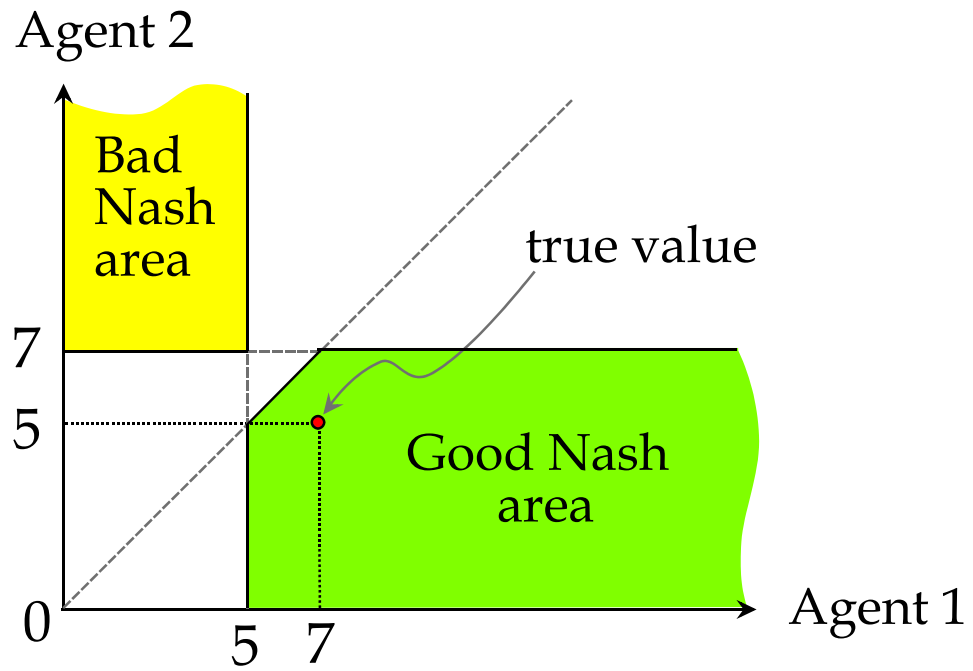


Figure 2: Equilibria of the Second Price Auction

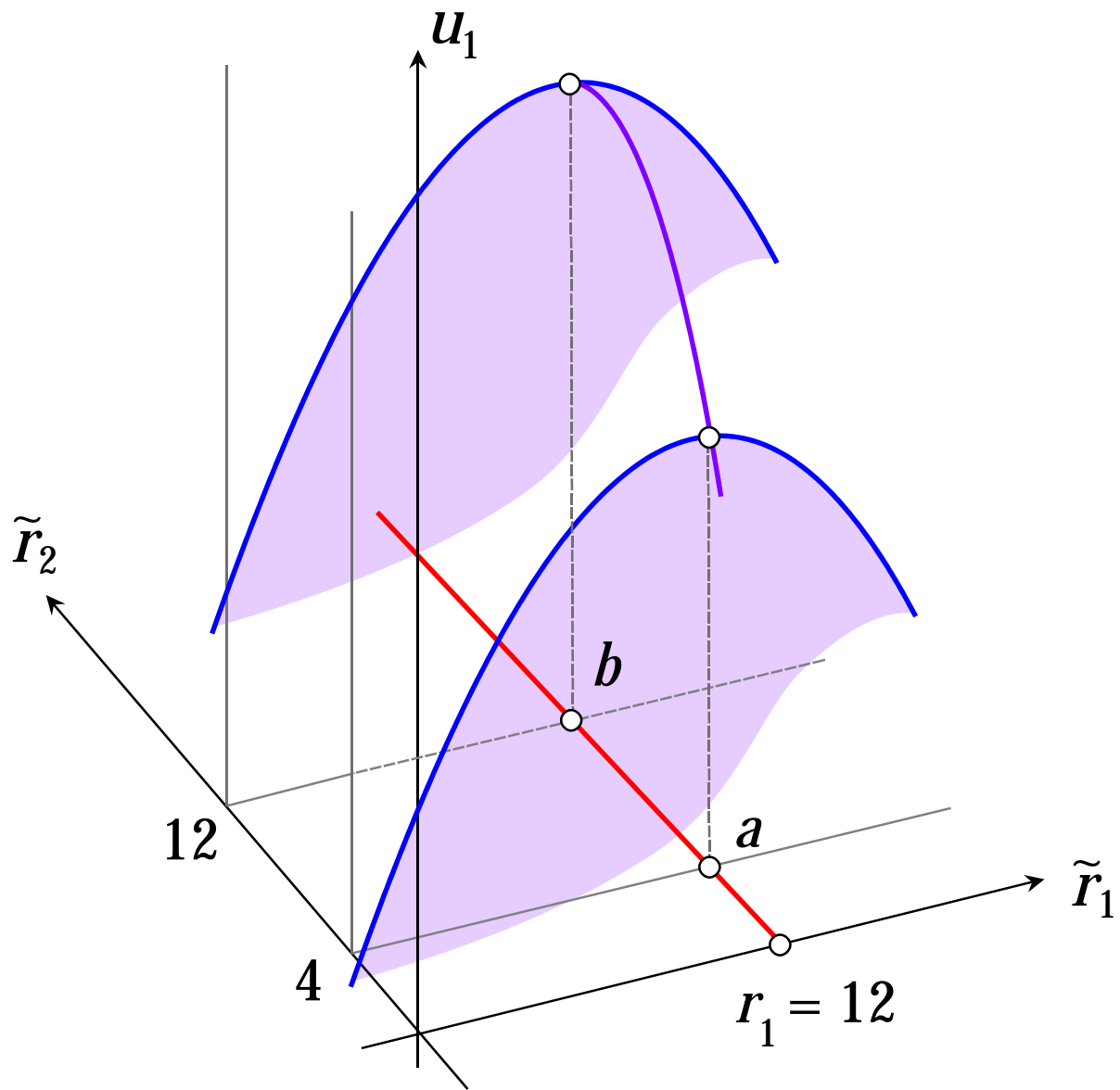


Figure 3: Payoff function of a Groves Mechanism with Single-Peaked Preferences

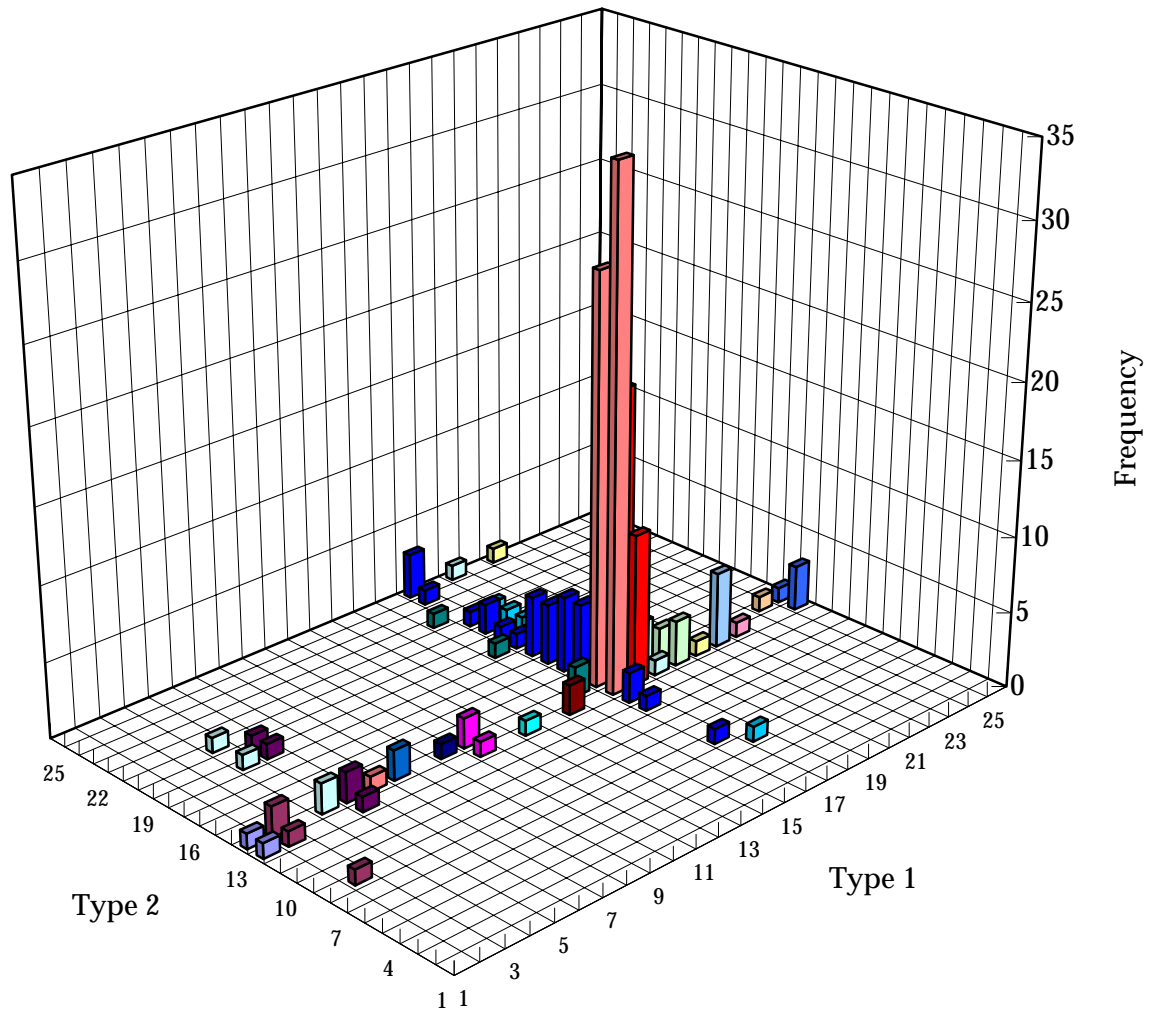


Figure 4: Treatment P -- All Pairs Choices

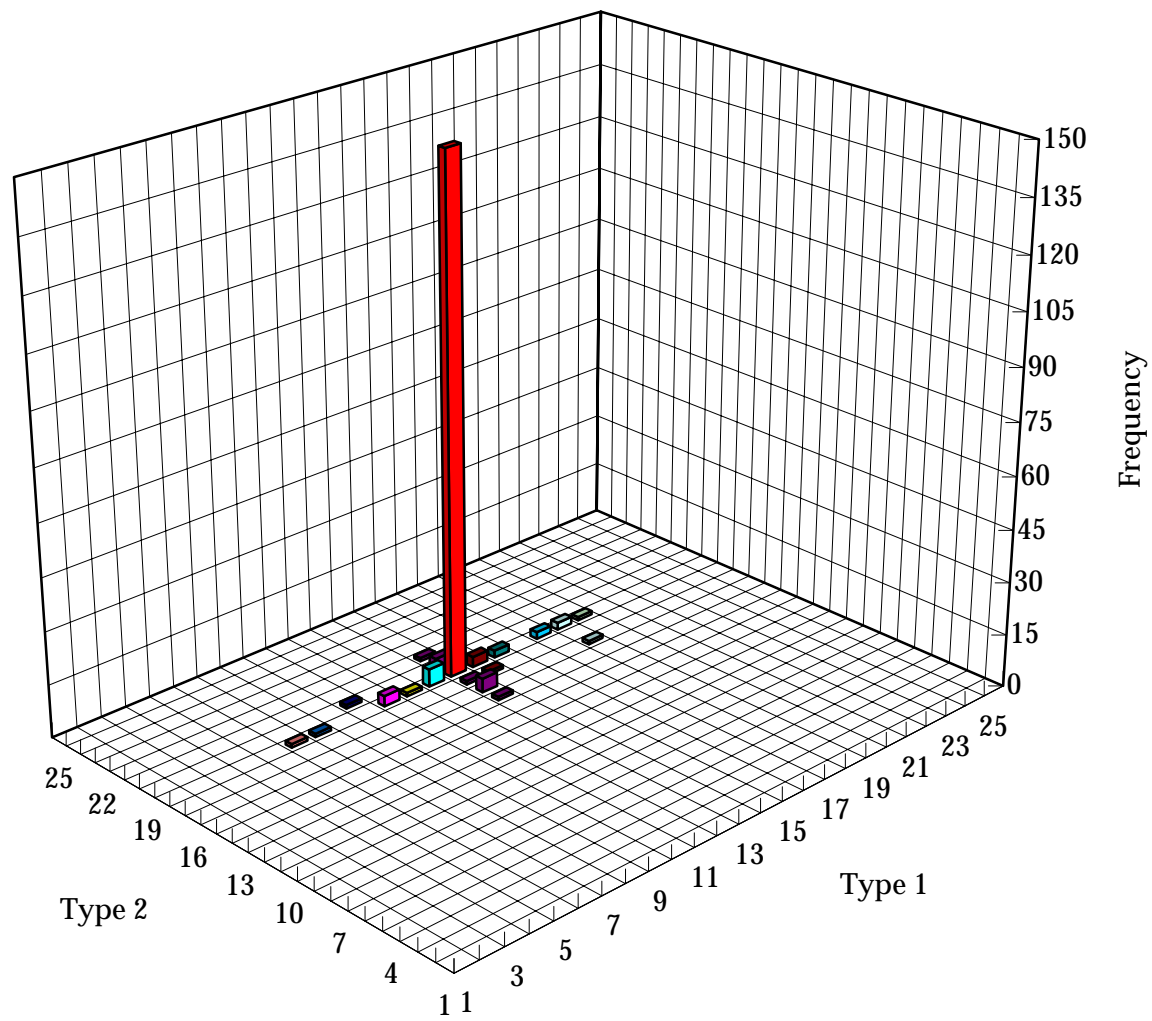


Figure 5: Treatment S -- All Pairs Choices

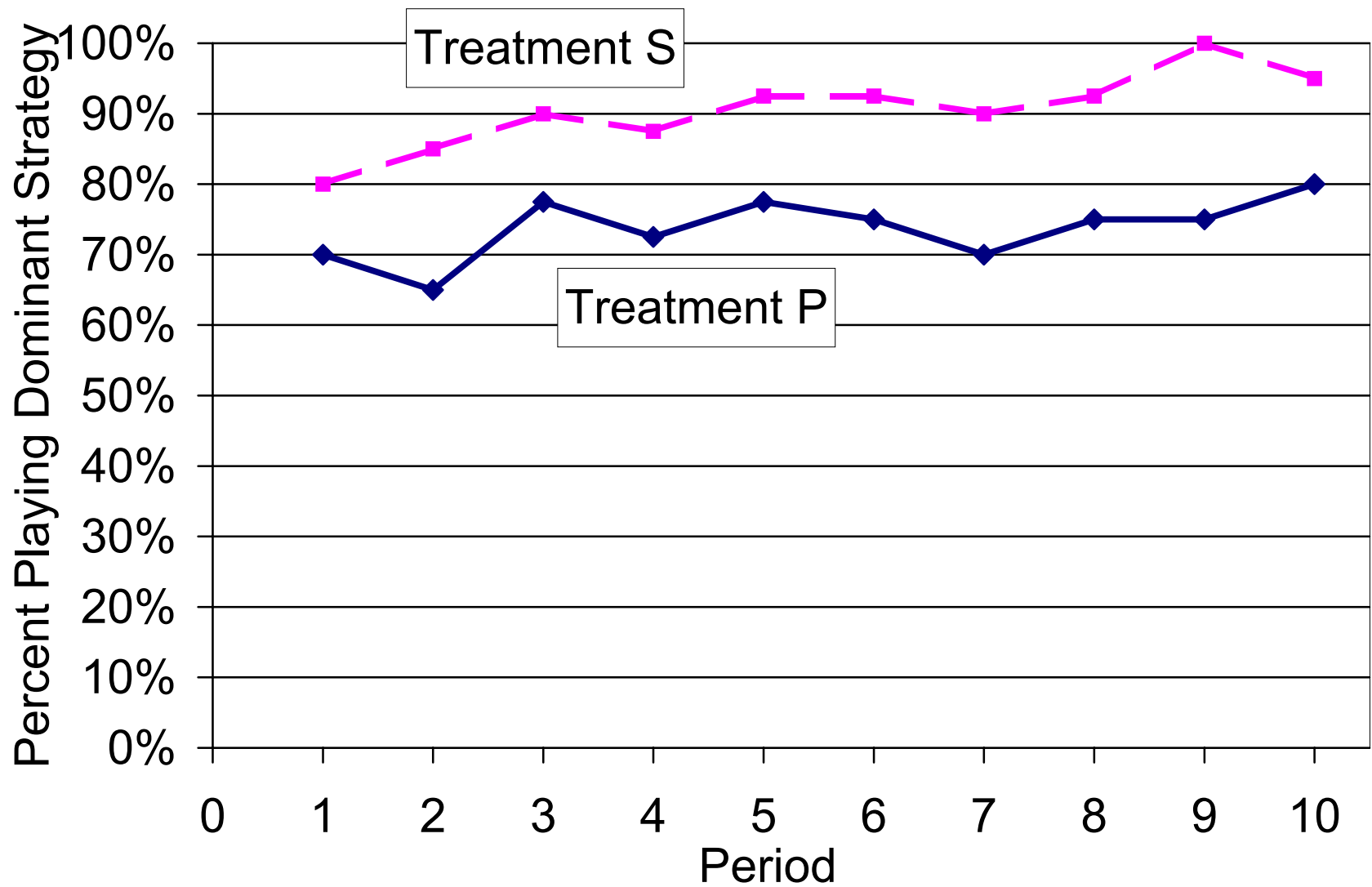


Figure 6: Rates that Individuals Play Dominant Strategies

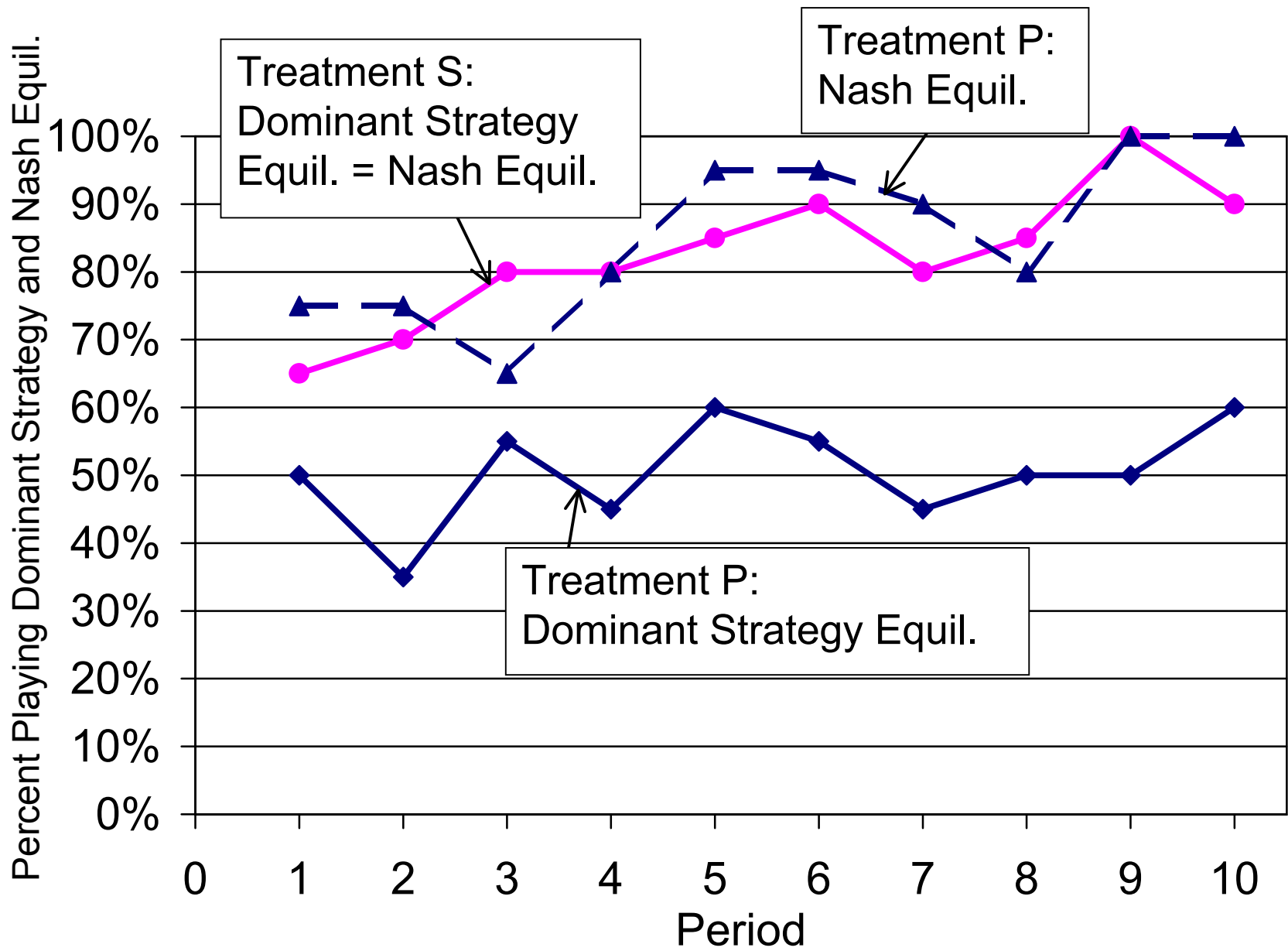


Figure 7: Rates that Pairs Play Dominant Strategy and Nash Equilibria



Dominant Strategy Equilibrium
 Pareto Dominated Dominant Strategy Equilibrium



Good Nash Equilibrium
 Bad Nash Equilibrium

P1

The number which you choose (Type 1)

The number
 which the
 other person
 chooses
 (Type 2)

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25
1	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294
2	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294
3	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	182
4	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	196	196
5	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	210	210	210
6	280	280	280	280	280	280	280	280	280	280	280	280	280	280	280	280	280	280	280	280	280	210	210	210	210
7	266	266	266	266	266	266	266	266	266	266	266	266	266	266	266	266	266	266	266	266	266	210	210	210	210
8	252	252	252	252	252	252	252	252	252	252	252	252	252	252	252	252	252	252	252	252	210	210	210	210	210
9	238	238	238	238	238	238	238	238	238	238	238	238	238	238	238	238	238	238	210	210	210	210	210	210	210
10	224	224	224	224	224	224	224	224	224	224	224	224	224	224	224	224	224	224	210	210	210	210	210	210	210
11	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210
12	196	196	196	196	196	196	196	196	196	196	196	196	196	196	196	196	210	210	210	210	210	210	210	210	210
13	182	182	182	182	182	182	182	182	182	182	182	182	182	182	210	210	210	210	210	210	210	210	210	210	210
14	168	168	168	168	168	168	168	168	168	168	168	168	168	210	210	210	210	210	210	210	210	210	210	210	210
15	154	154	154	154	154	154	154	154	154	154	154	154	210	210	210	210	210	210	210	210	210	210	210	210	210
16	140	140	140	140	140	140	140	140	140	140	140	210	210	210	210	210	210	210	210	210	210	210	210	210	210
17	126	126	126	126	126	126	126	126	126	126	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210
18	112	112	112	112	112	112	112	112	112	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210
19	98	98	98	98	98	98	98	98	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210
20	84	84	84	84	84	84	84	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210
21	70	70	70	70	70	70	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210
22	56	56	56	56	56	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210
23	42	42	42	42	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210
24	28	28	28	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210
25	14	14	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210

Table 1. Dominant Strategy Equilibria and Nash Equilibria for Payoff Table of Type 1 in Treatment P.



Dominant Strategy Equilibrium
 Pareto Dominated Dominant Strategy Equilibrium



Good Nash Equilibrium
 Bad Nash Equilibrium

P2

The number which you choose (Type 2)

The number which the other person chooses (Type 1)

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	
1	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182
2	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182
3	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	14
4	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	28	28
5	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	42	42	42
6	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	56	56	56	56
7	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	70	70	70	70	70
8	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	84	84	84	84	84	84
9	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	98	98	98	98	98	98	98
10	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	112	112	112	112	112	112	112
11	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	126	126	126	126	126	126	126	126	126	126
12	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	140	140	140	140	140	140	140	140	140	140	140
13	182	182	182	182	182	182	182	182	182	182	182	182	182	182	154	154	154	154	154	154	154	154	154	154	154	154
14	182	182	182	182	182	182	182	182	182	182	182	182	182	182	168	168	168	168	168	168	168	168	168	168	168	168
15	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182	182
16	182	182	182	182	182	182	182	182	182	182	182	182	196	196	196	196	196	196	196	196	196	196	196	196	196	196
17	182	182	182	182	182	182	182	182	182	182	182	210	210	210	210	210	210	210	210	210	210	210	210	210	210	210
18	182	182	182	182	182	182	182	182	182	224	224	224	224	224	224	224	224	224	224	224	224	224	224	224	224	224
19	182	182	182	182	182	182	182	182	238	238	238	238	238	238	238	238	238	238	238	238	238	238	238	238	238	238
20	182	182	182	182	182	182	182	252	252	252	252	252	252	252	252	252	252	252	252	252	252	252	252	252	252	252
21	182	182	182	182	182	182	266	266	266	266	266	266	266	266	266	266	266	266	266	266	266	266	266	266	266	266
22	182	182	182	182	182	280	280	280	280	280	280	280	280	280	280	280	280	280	280	280	280	280	280	280	280	280
23	182	182	182	182	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294
24	168	168	168	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294
25	154	154	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294	294

Table 2. Dominant Strategy Equilibria and Nash Equilibria for Payoff Table of Type 2 in Treatment P.

Payoff Table (for the Actual Experiment)

		The number which you choose																								
		1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25
The number which the other person chooses	1	116	124	131	138	144	150	154	158	161	164	166	167	167	167	166	164	161	158	154	150	144	138	131	124	116
	2	124	132	140	146	152	158	162	166	170	172	174	175	175	175	174	172	170	166	162	158	152	146	140	132	124
	3	131	140	147	154	160	165	170	174	177	180	181	182	183	182	181	180	177	174	170	165	160	154	147	140	131
	4	138	146	154	161	167	172	177	181	184	186	188	189	190	189	188	186	184	181	177	172	167	161	154	146	138
	5	144	152	160	167	173	178	183	187	190	192	194	195	196	195	194	192	190	187	183	178	173	167	160	152	144
	6	150	158	165	172	178	184	188	192	195	198	200	201	201	201	200	198	195	192	188	184	178	172	165	158	150
	7	154	162	170	177	183	188	193	197	200	202	204	205	206	205	204	202	200	197	193	188	183	177	170	162	154
	8	158	166	174	181	187	192	197	201	204	206	208	209	210	209	208	206	204	201	197	192	187	181	174	166	158
	9	161	170	177	184	190	195	200	204	207	210	211	212	213	212	211	210	207	204	200	195	190	184	177	170	161
	10	164	172	180	186	192	198	202	206	210	212	214	215	215	215	214	212	210	206	202	198	192	186	180	172	164
	11	166	174	181	188	194	200	204	208	211	214	216	217	217	217	216	214	211	208	204	200	194	188	181	174	166
	12	167	175	182	189	195	201	205	209	212	215	217	218	218	218	217	215	212	209	205	201	195	189	182	175	167
	13	167	175	183	190	196	201	206	210	213	215	217	218	219	218	217	215	213	210	206	201	196	190	183	175	167
	14	167	175	182	189	195	201	205	209	212	215	217	218	218	218	217	215	212	209	205	201	195	189	182	175	167
	15	166	174	181	188	194	200	204	208	211	214	216	217	217	217	216	214	211	208	204	200	194	188	181	174	166
	16	164	172	180	186	192	198	202	206	210	212	214	215	215	215	214	212	210	206	202	198	192	186	180	172	164
	17	161	170	177	184	190	195	200	204	207	210	211	212	213	212	211	210	207	204	200	195	190	184	177	170	161
	18	158	166	174	181	187	192	197	201	204	206	208	209	210	209	208	206	204	201	197	192	187	181	174	166	158
	19	154	162	170	177	183	188	193	197	200	202	204	205	206	205	204	202	200	197	193	188	183	177	170	162	154
	20	150	158	165	172	178	184	188	192	195	198	200	201	201	201	200	198	195	192	188	184	178	172	165	158	150
	21	144	152	160	167	173	178	183	187	190	192	194	195	196	195	194	192	190	187	183	178	173	167	160	152	144
	22	138	146	154	161	167	172	177	181	184	186	188	189	190	189	188	186	184	181	177	172	167	161	154	146	138
	23	131	140	147	154	160	165	170	174	177	180	181	182	183	182	181	180	177	174	170	165	160	154	147	140	131
	24	124	132	140	146	152	158	162	166	170	172	174	175	175	175	174	172	170	166	162	158	152	146	140	132	124
	25	116	124	131	138	144	150	154	158	161	164	166	167	167	167	166	164	161	158	154	150	144	138	131	124	116

Table 3. Payoff Table of Type 1 distributed in Treatment S.

Payoff Table (for the Actual Experiment)

		The number which you choose																								
		1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25
The number which the other person chooses	1	12	24	35	45	55	64	72	80	86	92	98	102	106	110	112	114	115	115	115	114	112	110	106	102	98
	2	24	36	47	57	67	76	84	91	98	104	110	114	118	121	124	126	127	127	127	126	124	121	118	114	110
	3	35	47	58	68	78	87	95	102	109	115	121	125	129	132	135	137	138	138	138	137	135	132	129	125	121
	4	45	57	68	79	88	97	105	113	120	126	131	136	140	143	145	147	148	149	148	147	145	143	140	136	131
	5	55	67	78	88	98	107	115	122	129	135	141	145	149	152	155	157	158	158	158	157	155	152	149	145	141
	6	64	76	87	97	107	116	124	131	138	144	150	154	158	161	164	166	167	167	167	166	164	161	158	154	150
	7	72	84	95	105	115	124	132	140	146	152	158	162	166	170	172	174	175	175	175	174	172	170	166	162	158
	8	80	91	102	113	122	131	140	147	154	160	165	170	174	177	180	181	182	183	182	181	180	177	174	170	165
	9	86	98	109	120	129	138	146	154	161	167	172	177	181	184	186	188	189	190	189	188	186	184	181	177	172
	10	92	104	115	126	135	144	152	160	167	173	178	183	187	190	192	194	195	196	195	194	192	190	187	183	178
	11	98	110	121	131	141	150	158	165	172	178	184	188	192	195	198	200	201	201	201	200	198	195	192	188	184
	12	102	114	125	136	145	154	162	170	177	183	188	193	197	200	202	204	205	206	205	204	202	200	197	193	188
	13	106	118	129	140	149	158	166	174	181	187	192	197	201	204	206	208	209	210	209	208	206	204	201	197	192
	14	110	121	132	143	152	161	170	177	184	190	195	200	204	207	210	211	212	213	212	211	210	207	204	200	195
	15	112	124	135	145	155	164	172	180	186	192	198	202	206	210	212	214	215	215	215	214	212	210	206	202	198
	16	114	126	137	147	157	166	174	181	188	194	200	204	208	211	214	216	217	217	217	216	214	211	208	204	200
	17	115	127	138	148	158	167	175	182	189	195	201	205	209	212	215	217	218	218	218	217	215	212	209	205	201
	18	115	127	138	149	158	167	175	183	190	196	201	206	210	213	215	217	218	219	218	217	215	213	210	206	201
	19	115	127	138	148	158	167	175	182	189	195	201	205	209	212	215	217	218	218	218	217	215	212	209	205	201
	20	114	126	137	147	157	166	174	181	188	194	200	204	208	211	214	216	217	217	217	216	214	211	208	204	200
	21	112	124	135	145	155	164	172	180	186	192	198	202	206	210	212	214	215	215	215	214	212	210	206	202	198
	22	110	121	132	143	152	161	170	177	184	190	195	200	204	207	210	211	212	213	212	211	210	207	204	200	195
	23	106	118	129	140	149	158	166	174	181	187	192	197	201	204	206	208	209	210	209	208	206	204	201	197	192
	24	102	114	125	136	145	154	162	170	177	183	188	193	197	200	202	204	205	206	205	204	202	200	197	193	188
	25	98	110	121	131	141	150	158	165	172	178	184	188	192	195	198	200	201	201	201	200	198	195	192	188	184

Table 4. Payoff Table of Type 2 distributed in Treatment S.

	(1) Individuals play dominant strategies	(2) Pairs play dominant strategy equilibrium
Dummy variable=1 for Treatment S	0.720** (0.346)	0.887** (0.143)
Dummy variable=1 for sessions at Purdue		0.170 (0.142)
Intercept	1.236** (0.266)	-0.095 (.0122)
$\rho = \sigma_u^2 / (\sigma_v^2 + \sigma_u^2)$ (random effects significance)	0.627** (0.069)	
Observations	720	360
Log-likelihood	-247.2	-211.3
Restricted log-likelihood	-344.5	-231.8

Notes: Standard errors shown in parentheses. ** denotes significantly different from zero at five-percent. Model in column (1) is estimated with a random subjects effect error term $u_i + v_{it}$.

Table 5. Probit Models of Individual and Pair Dominant Strategy Play

Appendix C: Payoff Tables for the Pivotal and Groves Mechanisms

\tilde{v}_1 type 1's reported value

\tilde{v}_2
type 2's
reported
value

	-22	-21	-20	-19	-18	-17	-16	-15	-14	-13	-12	-11	-10	-9	-8	-7	-6	-5	-4	-3	-2	-1	0	1	2
-4	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
-3	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
-2	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	-8
-1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	-7	-7
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	-6	-6	-6
1	-1	-1	-1	-1	-1	-1	-1	-1	-1	-1	-1	-1	-1	-1	-1	-1	-1	-1	-1	-1	-1	-6	-6	-6	-6
2	-2	-2	-2	-2	-2	-2	-2	-2	-2	-2	-2	-2	-2	-2	-2	-2	-2	-2	-2	-2	-6	-6	-6	-6	-6
3	-3	-3	-3	-3	-3	-3	-3	-3	-3	-3	-3	-3	-3	-3	-3	-3	-3	-3	-3	-6	-6	-6	-6	-6	-6
4	-4	-4	-4	-4	-4	-4	-4	-4	-4	-4	-4	-4	-4	-4	-4	-4	-4	-6	-6	-6	-6	-6	-6	-6	-6
5	-5	-5	-5	-5	-5	-5	-5	-5	-5	-5	-5	-5	-5	-5	-5	-5	-6	-6	-6	-6	-6	-6	-6	-6	-6
6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6
7	-7	-7	-7	-7	-7	-7	-7	-7	-7	-7	-7	-7	-7	-7	-7	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6
8	-8	-8	-8	-8	-8	-8	-8	-8	-8	-8	-8	-8	-8	-8	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6
9	-9	-9	-9	-9	-9	-9	-9	-9	-9	-9	-9	-9	-9	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6
10	-10	-10	-10	-10	-10	-10	-10	-10	-10	-10	-10	-10	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6
11	-11	-11	-11	-11	-11	-11	-11	-11	-11	-11	-11	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6
12	-12	-12	-12	-12	-12	-12	-12	-12	-12	-12	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6
13	-13	-13	-13	-13	-13	-13	-13	-13	-13	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6
14	-14	-14	-14	-14	-14	-14	-14	-14	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6
15	-15	-15	-15	-15	-15	-15	-15	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6
16	-16	-16	-16	-16	-16	-16	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6
17	-17	-17	-17	-17	-17	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6
18	-18	-18	-18	-18	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6
19	-19	-19	-19	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6
20	-20	-20	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6	-6

Table C-1. Payoff Table of Type 1 for the Pivotal Mechanism.

\bar{v}_2 type 2's reported value

\bar{v}_1
type 1's
reported
value

	-4	-3	-2	-1	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20		
-22	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0		
-21	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0		
-20	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	-12		
-19	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	-11	-11	
-18	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	-10	-10	-10	
-17	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	-9	-9	-9	-9
-16	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	-8	-8	-8	-8	-8
-15	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	-7	-7	-7	-7	-7	-7	-7
-14	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	-6	-6	-6	-6	-6	-6	-6	-6
-13	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	-5	-5	-5	-5	-5	-5	-5	-5	-5	-5
-12	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	-4	-4	-4	-4	-4	-4	-4	-4	-4	-4	-4
-11	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	-3	-3	-3	-3	-3	-3	-3	-3	-3	-3	-3	-3
-10	0	0	0	0	0	0	0	0	0	0	0	0	0	0	-2	-2	-2	-2	-2	-2	-2	-2	-2	-2	-2	-2	-2
-9	0	0	0	0	0	0	0	0	0	0	0	0	0	-1	-1	-1	-1	-1	-1	-1	-1	-1	-1	-1	-1	-1	-1
-8	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
-7	0	0	0	0	0	0	0	0	0	0	0	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1
-6	0	0	0	0	0	0	0	0	0	0	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2
-5	0	0	0	0	0	0	0	0	0	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3	3
-4	0	0	0	0	0	0	0	0	4	4	4	4	4	4	4	4	4	4	4	4	4	4	4	4	4	4	4
-3	0	0	0	0	0	0	0	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5
-2	0	0	0	0	0	0	6	6	6	6	6	6	6	6	6	6	6	6	6	6	6	6	6	6	6	6	6
-1	0	0	0	0	0	7	7	7	7	7	7	7	7	7	7	7	7	7	7	7	7	7	7	7	7	7	7
0	0	0	0	0	8	8	8	8	8	8	8	8	8	8	8	8	8	8	8	8	8	8	8	8	8	8	8
1	-1	-1	-1	8	8	8	8	8	8	8	8	8	8	8	8	8	8	8	8	8	8	8	8	8	8	8	8
2	-2	-2	8	8	8	8	8	8	8	8	8	8	8	8	8	8	8	8	8	8	8	8	8	8	8	8	8

Table C-2. Payoff Table of Type 2 for the Pivotal Mechanism.

$\bar{\tau}_1$ type 1's reported peak

		0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
type 2's reported peak	0	-144.0	-132.5	-122.0	-112.5	-104.0	-96.5	-90.0	-84.5	-80.0	-76.5	-74.0	-72.5	-72.0	-72.5	-74.0	-76.5	-80.0	-84.5	-90.0	-96.5	-104.0	-112.5	-122.0	-132.5	-144.0
	1	-132.5	-121.0	-110.5	-101.0	-92.5	-85.0	-78.5	-73.0	-68.5	-65.0	-62.5	-61.0	-60.5	-61.0	-62.5	-65.0	-68.5	-73.0	-78.5	-85.0	-92.5	-101.0	-110.5	-121.0	-132.5
	2	-122.0	-110.5	-100.0	-90.5	-82.0	-74.5	-68.0	-62.5	-58.0	-54.5	-52.0	-50.5	-50.0	-50.5	-52.0	-54.5	-58.0	-62.5	-68.0	-74.5	-82.0	-90.5	-100.0	-110.5	-122.0
	3	-112.5	-101.0	-90.5	-81.0	-72.5	-65.0	-58.5	-53.0	-48.5	-45.0	-42.5	-41.0	-40.5	-41.0	-42.5	-45.0	-48.5	-53.0	-58.5	-65.0	-72.5	-81.0	-90.5	-101.0	-112.5
	4	-104.0	-92.5	-82.0	-72.5	-64.0	-56.5	-50.0	-44.5	-40.0	-36.5	-34.0	-32.5	-32.0	-32.5	-34.0	-36.5	-40.0	-44.5	-50.0	-56.5	-64.0	-72.5	-82.0	-92.5	-104.0
	5	-96.5	-85.0	-74.5	-65.0	-56.5	-49.0	-42.5	-37.0	-32.5	-29.0	-26.5	-25.0	-24.5	-25.0	-26.5	-29.0	-32.5	-37.0	-42.5	-49.0	-56.5	-65.0	-74.5	-85.0	-96.5
	6	-90.0	-78.5	-68.0	-58.5	-50.0	-42.5	-36.0	-30.5	-26.0	-22.5	-20.0	-18.5	-18.0	-18.5	-20.0	-22.5	-26.0	-30.5	-36.0	-42.5	-50.0	-58.5	-68.0	-78.5	-90.0
	7	-84.5	-73.0	-62.5	-53.0	-44.5	-37.0	-30.5	-25.0	-20.5	-17.0	-14.5	-13.0	-12.5	-13.0	-14.5	-17.0	-20.5	-25.0	-30.5	-37.0	-44.5	-53.0	-62.5	-73.0	-84.5
	8	-80.0	-68.5	-58.0	-48.5	-40.0	-32.5	-26.0	-20.5	-16.0	-12.5	-10.0	-8.5	-8.0	-8.5	-10.0	-12.5	-16.0	-20.5	-26.0	-32.5	-40.0	-48.5	-58.0	-68.5	-80.0
	9	-76.5	-65.0	-54.5	-45.0	-36.5	-29.0	-22.5	-17.0	-12.5	-9.0	-6.5	-5.0	-4.5	-5.0	-6.5	-9.0	-12.5	-17.0	-22.5	-29.0	-36.5	-45.0	-54.5	-65.0	-76.5
	10	-74.0	-62.5	-52.0	-42.5	-34.0	-26.5	-20.0	-14.5	-10.0	-6.5	-4.0	-2.5	-2.0	-2.5	-4.0	-6.5	-10.0	-14.5	-20.0	-26.5	-34.0	-42.5	-52.0	-62.5	-74.0
	11	-72.5	-61.0	-50.5	-41.0	-32.5	-25.0	-18.5	-13.0	-8.5	-5.0	-2.5	-1.0	-0.5	-1.0	-2.5	-5.0	-8.5	-13.0	-18.5	-25.0	-32.5	-41.0	-50.5	-61.0	-72.5
	12	-72.0	-60.5	-50.0	-40.5	-32.0	-24.5	-18.0	-12.5	-8.0	-4.5	-2.0	-0.5	0.0	-0.5	-2.0	-4.5	-8.0	-12.5	-18.0	-24.5	-32.0	-40.5	-50.0	-60.5	-72.0
	13	-72.5	-61.0	-50.5	-41.0	-32.5	-25.0	-18.5	-13.0	-8.5	-5.0	-2.5	-1.0	-0.5	-1.0	-2.5	-5.0	-8.5	-13.0	-18.5	-25.0	-32.5	-41.0	-50.5	-61.0	-72.5
	14	-74.0	-62.5	-52.0	-42.5	-34.0	-26.5	-20.0	-14.5	-10.0	-6.5	-4.0	-2.5	-2.0	-2.5	-4.0	-6.5	-10.0	-14.5	-20.0	-26.5	-34.0	-42.5	-52.0	-62.5	-74.0
	15	-76.5	-65.0	-54.5	-45.0	-36.5	-29.0	-22.5	-17.0	-12.5	-9.0	-6.5	-5.0	-4.5	-5.0	-6.5	-9.0	-12.5	-17.0	-22.5	-29.0	-36.5	-45.0	-54.5	-65.0	-76.5
	16	-80.0	-68.5	-58.0	-48.5	-40.0	-32.5	-26.0	-20.5	-16.0	-12.5	-10.0	-8.5	-8.0	-8.5	-10.0	-12.5	-16.0	-20.5	-26.0	-32.5	-40.0	-48.5	-58.0	-68.5	-80.0
	17	-84.5	-73.0	-62.5	-53.0	-44.5	-37.0	-30.5	-25.0	-20.5	-17.0	-14.5	-13.0	-12.5	-13.0	-14.5	-17.0	-20.5	-25.0	-30.5	-37.0	-44.5	-53.0	-62.5	-73.0	-84.5
	18	-90.0	-78.5	-68.0	-58.5	-50.0	-42.5	-36.0	-30.5	-26.0	-22.5	-20.0	-18.5	-18.0	-18.5	-20.0	-22.5	-26.0	-30.5	-36.0	-42.5	-50.0	-58.5	-68.0	-78.5	-90.0
	19	-96.5	-85.0	-74.5	-65.0	-56.5	-49.0	-42.5	-37.0	-32.5	-29.0	-26.5	-25.0	-24.5	-25.0	-26.5	-29.0	-32.5	-37.0	-42.5	-49.0	-56.5	-65.0	-74.5	-85.0	-96.5
	20	-104.0	-92.5	-82.0	-72.5	-64.0	-56.5	-50.0	-44.5	-40.0	-36.5	-34.0	-32.5	-32.0	-32.5	-34.0	-36.5	-40.0	-44.5	-50.0	-56.5	-64.0	-72.5	-82.0	-92.5	-104.0
	21	-112.5	-101.0	-90.5	-81.0	-72.5	-65.0	-58.5	-53.0	-48.5	-45.0	-42.5	-41.0	-40.5	-41.0	-42.5	-45.0	-48.5	-53.0	-58.5	-65.0	-72.5	-81.0	-90.5	-101.0	-112.5
	22	-122.0	-110.5	-100.0	-90.5	-82.0	-74.5	-68.0	-62.5	-58.0	-54.5	-52.0	-50.5	-50.0	-50.5	-52.0	-54.5	-58.0	-62.5	-68.0	-74.5	-82.0	-90.5	-100.0	-110.5	-122.0
	23	-132.5	-121.0	-110.5	-101.0	-92.5	-85.0	-78.5	-73.0	-68.5	-65.0	-62.5	-61.0	-60.5	-61.0	-62.5	-65.0	-68.5	-73.0	-78.5	-85.0	-92.5	-101.0	-110.5	-121.0	-132.5
	24	-144.0	-132.5	-122.0	-112.5	-104.0	-96.5	-90.0	-84.5	-80.0	-76.5	-74.0	-72.5	-72.0	-72.5	-74.0	-76.5	-80.0	-84.5	-90.0	-96.5	-104.0	-112.5	-122.0	-132.5	-144.0

Table C-3. Payoff Table of Type 1 for a Groves Mechanism with Single-Peaked Preferences.

\bar{v}_2 type 2's reported peak

		0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	
type 1's reported peak	\bar{v}_1	0	-289.0	-272.5	-257.0	-242.5	-229.0	-216.5	-205.0	-194.5	-185.0	-176.5	-169.0	-162.5	-157.0	-152.5	-149.0	-146.5	-145.0	-144.5	-145.0	-146.5	-149.0	-152.5	-157.0	-162.5	-169.0
	1	-272.5	-256.0	-240.5	-226.0	-212.5	-200.0	-188.5	-178.0	-168.5	-160.0	-152.5	-146.0	-140.5	-136.0	-132.5	-130.0	-128.5	-128.0	-128.5	-130.0	-132.5	-136.0	-140.5	-146.0	-152.5	
	2	-257.0	-240.5	-225.0	-210.5	-197.0	-184.5	-173.0	-162.5	-153.0	-144.5	-137.0	-130.5	-125.0	-120.5	-117.0	-114.5	-113.0	-112.5	-113.0	-114.5	-117.0	-120.5	-125.0	-130.5	-137.0	
	3	-242.5	-226.0	-210.5	-196.0	-182.5	-170.0	-158.5	-148.0	-138.5	-130.0	-122.5	-116.0	-110.5	-106.0	-102.5	-100.0	-98.5	-98.0	-98.5	-100.0	-102.5	-106.0	-110.5	-116.0	-122.5	
	4	-229.0	-212.5	-197.0	-182.5	-169.0	-156.5	-145.0	-134.5	-125.0	-116.5	-109.0	-102.5	-97.0	-92.5	-89.0	-86.5	-85.0	-84.5	-85.0	-86.5	-89.0	-92.5	-97.0	-102.5	-109.0	
	5	-216.5	-200.0	-184.5	-170.0	-156.5	-144.0	-132.5	-122.0	-112.5	-104.0	-96.5	-90.0	-84.5	-80.0	-76.5	-74.0	-72.5	-72.0	-72.5	-74.0	-76.5	-80.0	-84.5	-90.0	-96.5	
	6	-205.0	-188.5	-173.0	-158.5	-145.0	-132.5	-121.0	-110.5	-101.0	-92.5	-85.0	-78.5	-73.0	-68.5	-65.0	-62.5	-61.0	-60.5	-61.0	-62.5	-65.0	-68.5	-73.0	-78.5	-85.0	
	7	-194.5	-178.0	-162.5	-148.0	-134.5	-122.0	-110.5	-100.0	-90.5	-82.0	-74.5	-68.0	-62.5	-58.0	-54.5	-52.0	-50.5	-50.0	-50.5	-52.0	-54.5	-58.0	-62.5	-68.0	-74.5	
	8	-185.0	-168.5	-153.0	-138.5	-125.0	-112.5	-101.0	-90.5	-81.0	-72.5	-65.0	-58.5	-53.0	-48.5	-45.0	-42.5	-41.0	-40.5	-41.0	-42.5	-45.0	-48.5	-53.0	-58.5	-65.0	
	9	-176.5	-160.0	-144.5	-130.0	-116.5	-104.0	-92.5	-82.0	-72.5	-64.0	-56.5	-50.0	-44.5	-40.0	-36.5	-34.0	-32.5	-32.0	-32.5	-34.0	-36.5	-40.0	-44.5	-50.0	-56.5	
	10	-169.0	-152.5	-137.0	-122.5	-109.0	-96.5	-85.0	-74.5	-65.0	-56.5	-49.0	-42.5	-37.0	-32.5	-29.0	-26.5	-25.0	-24.5	-25.0	-26.5	-29.0	-32.5	-37.0	-42.5	-49.0	
	11	-162.5	-146.0	-130.5	-116.0	-102.5	-90.0	-78.5	-68.0	-58.5	-50.0	-42.5	-36.0	-30.5	-26.0	-22.5	-20.0	-18.5	-18.0	-18.5	-20.0	-22.5	-26.0	-30.5	-36.0	-42.5	
	12	-157.0	-140.5	-125.0	-110.5	-97.0	-84.5	-73.0	-62.5	-53.0	-44.5	-37.0	-30.5	-25.0	-20.5	-17.0	-14.5	-13.0	-12.5	-13.0	-14.5	-17.0	-20.5	-25.0	-30.5	-37.0	
	13	-152.5	-136.0	-120.5	-106.0	-92.5	-80.0	-68.5	-58.0	-48.5	-40.0	-32.5	-26.0	-20.5	-16.0	-12.5	-10.0	-8.5	-8.0	-8.5	-10.0	-12.5	-16.0	-20.5	-26.0	-32.5	
	14	-149.0	-132.5	-117.0	-102.5	-89.0	-76.5	-65.0	-54.5	-45.0	-36.5	-29.0	-22.5	-17.0	-12.5	-9.0	-6.5	-5.0	-4.5	-5.0	-6.5	-9.0	-12.5	-17.0	-22.5	-29.0	
	15	-146.5	-130.0	-114.5	-100.0	-86.5	-74.0	-62.5	-52.0	-42.5	-34.0	-26.5	-20.0	-14.5	-10.0	-6.5	-4.0	-2.5	-2.0	-2.5	-4.0	-6.5	-10.0	-14.5	-20.0	-26.5	
	16	-145.0	-128.5	-113.0	-98.5	-85.0	-72.5	-61.0	-50.5	-41.0	-32.5	-25.0	-18.5	-13.0	-8.5	-5.0	-2.5	-1.0	-0.5	-1.0	-2.5	-5.0	-8.5	-13.0	-18.5	-25.0	
	17	-144.5	-128.0	-112.5	-98.0	-84.5	-72.0	-60.5	-50.0	-40.5	-32.0	-24.5	-18.0	-12.5	-8.0	-4.5	-2.0	-0.5	0.0	-0.5	-2.0	-4.5	-8.0	-12.5	-18.0	-24.5	
	18	-145.0	-128.5	-113.0	-98.5	-85.0	-72.5	-61.0	-50.5	-41.0	-32.5	-25.0	-18.5	-13.0	-8.5	-5.0	-2.5	-1.0	-0.5	-1.0	-2.5	-5.0	-8.5	-13.0	-18.5	-25.0	
	19	-146.5	-130.0	-114.5	-100.0	-86.5	-74.0	-62.5	-52.0	-42.5	-34.0	-26.5	-20.0	-14.5	-10.0	-6.5	-4.0	-2.5	-2.0	-2.5	-4.0	-6.5	-10.0	-14.5	-20.0	-26.5	
	20	-149.0	-132.5	-117.0	-102.5	-89.0	-76.5	-65.0	-54.5	-45.0	-36.5	-29.0	-22.5	-17.0	-12.5	-9.0	-6.5	-5.0	-4.5	-5.0	-6.5	-9.0	-12.5	-17.0	-22.5	-29.0	
	21	-152.5	-136.0	-120.5	-106.0	-92.5	-80.0	-68.5	-58.0	-48.5	-40.0	-32.5	-26.0	-20.5	-16.0	-12.5	-10.0	-8.5	-8.0	-8.5	-10.0	-12.5	-16.0	-20.5	-26.0	-32.5	
	22	-157.0	-140.5	-125.0	-110.5	-97.0	-84.5	-73.0	-62.5	-53.0	-44.5	-37.0	-30.5	-25.0	-20.5	-17.0	-14.5	-13.0	-12.5	-13.0	-14.5	-17.0	-20.5	-25.0	-30.5	-37.0	
	23	-162.5	-146.0	-130.5	-116.0	-102.5	-90.0	-78.5	-68.0	-58.5	-50.0	-42.5	-36.0	-30.5	-26.0	-22.5	-20.0	-18.5	-18.0	-18.5	-20.0	-22.5	-26.0	-30.5	-36.0	-42.5	
	24	-169.0	-152.5	-137.0	-122.5	-109.0	-96.5	-85.0	-74.5	-65.0	-56.5	-49.0	-42.5	-37.0	-32.5	-29.0	-26.5	-25.0	-24.5	-25.0	-26.5	-29.0	-32.5	-37.0	-42.5	-49.0	

Table C-4. Payoff Table of Type 2 for a Groves Mechanism with Single-Peaked Preferences.