# ON HANDOFF PERFORMANCE FOR AN INTEGRATED VOICE/DATA CELLULAR SYSTEM

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Abstract - One of the key performance in wireless cellular system is the handoff dropping probability. This paper studies the handoff performance for an integrated voice/data wireless cellular system. Specifically, we propose a new handoff control policy, called Dual Trunk Reservation with Queuing (DTR-Q). It uses two thresholds, one to reserve bandwidth for voice handoff, and the other for managing the data traffic. By taking advantage of the adaptiveness of data traffic, it uses a single queue to further buffer the data request in case the requested bandwidth is not available. The unique features of the proposed DTR-Q scheme are: 1) it can provide the necessary service guarantee and service differentiation for voice and data traffic; 2) It adopts a Complete Sharing approach that can maximize the channel utilization. We propose an analytical model to calculate the key performance measures, and thoroughly investigate system performances under a variety of system parameters.

**Keywords** – bandwidth allocation, guarded channel policy, handoff, mobile cellular networks

## 1. INTRODUCTION

We have recently witnessed phenomenal growth in the development and deployment of wireless services, evident from the proliferation of cellular data services and the emerging wireless multimedia applications. This opens up a new research avenue and calls for the re-examination of some of the fundamental issues in wireless cellular networks, in particular various issues such as QoS guarantee have to be carefully examined. One of the key elements in providing QoS guarantee is an effective bandwidth allocation policy, which not only has to ensure that the system guarantee potentially different QoS requirements and provide the necessary service differentiation from diverse applications, but at the same time has to fully utilize the scarce wireless bandwidth available in the wireless cellular networks. Moving from a voice-centric macro-cell wireless network to a multi-services micro-cell or/and pico-cell wireless network brings several new challenges: 1) the characteristics of data and multimedia applications are typically different from that of voice, e.g., data traffic usually consumes more bandwidths that of voice, but can be adaptive, thus has less stringent service requirement; 2) The limited wireless resources have to be utilized effectively, and fairly allocated to different types of traffic with

potentially different service requirements and service differentiation; 3) smaller cell potentially results in more handoffs, making it more difficult to provide the necessary QoS guarantee.

Bandwidth allocation has been extensively studied in single-service wireless cellular networks. The Guarded Channel (GC) schemes have been shown to be effective for providing the necessary QoS guarantee in terms of both call termination and call blocking probabilities [1-3]. One of the challenges in moving to a multi-service system is that the limited bandwidth has to be shared among multiple traffics. In [4] a *Complete Sharing (CS)* and *Complete Partition (CP)* schemes were in investigated for two types of traffic, namely narrow-band and wideband. It assumed that wideband traffic does not have handoff. Huang et al proposed a *movable boundary* bandwidth allocation scheme for voice and data traffic [5], in which bandwidth is divided into two sub-pools by two thresholds that can be dynamically adjusted. This facilitates the bandwidth provisioning for different QoS requirement and is adaptive to the changing traffic. The limitation is that there is no service differentiation between voice new calls and handoff. On the other hand, queuing the new or handoff calls can further reduce the handoff dropping probability as indicated in [1, 6].

In this paper, we propose a new handoff control scheme called *Dual Trunk Reservation with Queuing (DTR-Q)* scheme for a voice/data integrated cellular system. It uses *two thresholds*, one to reserve bandwidth for handoff voice calls, and the other for managing data traffics. By taking advantage of the adaptiveness of the data traffic, a single queue is used to further buffer the data requests in case the requested bandwidth is not available. The unique features of the proposed *DTR-Q* scheme are: 1) it can provide the necessary service guarantees and service differentiation for voice and data traffic; 2) It adopts the scheme of *Complete Sharing*, which can maximize the channel utilization. We propose an analytical model to calculate the key performance measures, and thoroughly investigate system performances under a variety of system parameters.

The rest of the paper is organized as follow. We describe the DTR-Q schemes in Section 2. We present the analytical models in Section 3, following by the numerical studies in Section 4. We conclude the paper in Section 5 with discussions on possible avenue for further study.

### 2. THE DTR-Q CONTROL POLICY

A good bandwidth allocation scheme should provide different *QoS* guarantees for different traffic types, while at the same time has to fully utilize the scarce wireless bandwidth available in the wireless cellular networks. A variety of bandwidth allocation schemes have been proposed to support multiple traffics, which can be classified as *Complete Partition (CP), Complete Sharing (CS)* or hybrid schemes, depending on how the bandwidth are allocating among diverse traffic. Compared with *CP* scheme, *CS* scheme can usually achieve much higher channel utilization [4]. This is because in *CS* scheme, users from different traffic types are allowed to share all the available channels statistically. Thus in this paper, we adopt *CS* control scheme in our *Dual Trunk Reservation with Queuing (DTR-Q)* scheme to maximize system efficiency.

In the DTR-Q scheme, total C channels in each cell are divided into *three* regions by two thresholds  $K_1$  and  $K_2$ (see Fig.1). When the number of channels occupied is less than the threshold  $K_2$ , then both data and voice traffic can be admitted into the system; when the number of channels occupied is over the threshold  $K_2$ , data calls will be placed into the queue if it is not full (assume the size of the finite queue is  $K_{\text{max}}$ ), otherwise it will be dropped; when the number of channels occupied is more than the threshold  $K_1$ , then only voice handoff calls can be allowed. The voice handoff call will be dropped only if there is no channel available. Under this basic control model, the voice handoff gets highest priority, while data receives lowest service. The reason is that the data traffic can tolerate certain degree of delay/delay-jitter, while voice cannot. On the other hand, buffering data requests can lead to a relative lower data blocking probability as shown in [6].

## 3. THE ANALYTICAL MODEL

We consider a homogeneous wireless network where all cells have the same number of channels and experience the same new and handoff call arrival rates. In each cell, the arrivals of new voice calls, new data calls, handoff voice calls and handoff data calls are Poisson distributed with arrival rate  $\lambda_{vn}$ ,  $\lambda_{dn}$ ,  $\lambda_{vh}$  and  $\lambda_{dh}$  respectively. Thus the total voice call arrival rate and data call arrival rate are  $\lambda_v = \lambda_{vn} + \lambda_{vh}$  and  $\lambda_d = \lambda_{dn} + \lambda_{dh}$  respectively. Since data can usually tolerate some degree of service degradation, new data calls and handoff data calls are not distinguished. Call duration times or call holding times of voice and data are exponentially distributed with the average call duration time  $1/\mu_{vr}$  and  $1/\mu_{dr}$ . In addition, the cell residence time for voice and data calls is exponentially distributed with mean  $1/\mu_{vh}$  and  $1/\mu_{dh}$  respectively. Thus, the channel occupancy times for voice and data calls are exponentially  $1/\mu_v = 1/(\mu_{vr} + \mu_{vh})$ distributed with mean and  $1/\mu_d = 1/(\mu_{dr} + \mu_{dh})$  respectively. This set of assumptions have been found to be reasonable as long as the number of mobiles in a cell is much greater than the number of channels and have been widely used in literature [7, 8, 2, 3]. The exponential channel holding time assumption is shown to be valid for a wide range of system under certain conditions [9], which allows us to conduct the mean value analysis [10]. We further assume that both voice and data occupy one unit of bandwidth.

Scheme DTR-Q can be modelled as a three-dimensional Markov chain. Let  $P_{i,j,k}$  be the steady state probability that there are *i* voice calls, *j* data calls in the system, and *k* data calls in the data buffer. The steady-state balance-equations of DTR-Q scheme are shown as below. Due to the length limitation, detail descriptions of the analytical models are given in [11].

Case 1: If 
$$i + j = 0$$
, then  
 $(\lambda_v + \lambda_d) P_{0,0,0} = \mu_v P_{1,0,0} + \mu_d P_{0,1,0}$  (1)

Case 2: If  $0 < i + j < K_2$ , then

$$(\lambda_{v} + \lambda_{d} + i\mu_{v} + j\mu_{d})P_{i,j,k} = (i+1)\mu_{v}P_{i+1,j,0} + (j+1)\mu_{d}P_{i,j+1,0} + \lambda_{v}P_{i-1,j,0} + \lambda_{d}P_{i,j-1,0}$$
(2)

Case 3: If 
$$i + j = K_2$$
 and  $k = 0$ , then  
 $(\lambda_{\nu} + \lambda_d + i\mu_{\nu} + j\mu_d)P_{i,j,0} = (i+1)\mu_{\nu}P_{i+1,j-1,1}$   
 $+ (j+1)\mu_d P_{i,j+1,0} + (i+1)\mu_{\nu}P_{i+1,j,0} + j\mu_d P_{i,j,1}$  (3)  
 $+ \lambda_{\nu}P_{i-1,j,0} + \lambda_d P_{i,j-1,0}$ 

Case 4: If 
$$i + j = K_2$$
 and  $k > 0$ , then  
 $(\lambda_v + \lambda_d + i\mu_v + j\mu_d)P_{i,j,k} = (i+1)\mu_v P_{i+1,j-1,k+1}$   
 $+ (j+1)\mu_d P_{i,j+1,k} + (i+1)\mu_v P_{i+1,j,k} + j\mu_d P_{i,j,k+1}$  (4)  
 $+ \lambda_d P_{i,j-1,k}$ 

Case 5: If 
$$K_2 < i + j < K_1$$
, then  
 $(\lambda_v + \lambda_d + i\mu_v + j\mu_d)P_{i,j,k} = (i+1)\mu_v P_{i+1,j1,k}$   
 $+ (j+1)\mu_d P_{i,j+1,k} + \lambda_v P_{i-1,j,k} + \lambda_d P_{i,j,k-1}$ 
(5)

Case 6: if 
$$i + j = K_1$$
, then  
 $(\lambda_v + \lambda_d + i\mu_v + j\mu_d)P_{i,j,k} = (i+1)\mu_v P_{i+1,j,k}$   
 $+ (j+1)\mu_d P_{i,j+1,k} + \lambda_v P_{i-1,j,k} + \lambda_d P_{i,j,k-1}$ 
(6)

Case 7: if 
$$K_1 < i + j < C$$
, then  
 $(\lambda_{vh} + \lambda_d + i\mu_v + j\mu_d)P_{i,j,k} = (i+1)\mu_v P_{i+1,j,k}$   
 $+ (j+1)\mu_d P_{i,j+1,k} + \lambda_{vh} P_{i-1,j,k} + \lambda_d P_{i,j,k-1}$ 
(7)

*Case 8: if* i + j = C, then

$$(\lambda_d + i\mu_v + j\mu_d)P_{i,j,k} = \lambda_{vh}P_{i-1,j,k} + \lambda_d P_{i,j,k-1}$$
(8)

After obtaining the steady-state transition equations for each state in the Markov chain, we can solve the linear equation together with the normalizing condition (see Eq. (9)) by using LU factorization.

$$\sum_{j=0}^{C} \sum_{k\geq 0}^{K_{\max}} P_{i,j,k} = 1$$
(9)

After obtaining all the steady state probabilities  $P_{i,j,k}$ , the voice call blocking probability  $P_{vb}$ , the handoff voice call dropping probability  $P_{vd}$ , the data call blocking probability  $P_{db}$  and the average queue length  $\overline{Q}$  can be derived as below:

$$P_{vb} = \sum_{i+j \ge K_1}^{C} \sum_{k \ge 0}^{K_{\text{max}}} P_{i,j,k}$$
(10)

$$P_{vd} = \sum_{i+j=C} \sum_{k\geq 0}^{K_{\max}} P_{i,j,k}$$
(11)

$$P_{db} = \sum_{i+j \ge K_2}^{C} \sum_{k \ge 0}^{K_{\text{max}}} P_{i,j,k}$$
(12)

$$\overline{Q} = \sum_{i+j=K_2}^{C} \sum_{k=K_{\text{max}}} k P_{i,j,k}$$
(13)

### 4. NUMERICAL RESULTS

In this section, we present numerical results in three cases: 1) performance under different voice call arrival rate  $\lambda_v$ ; 2) performance under different data call arrival rate  $\lambda_d$ ; 3) performance under different  $K_2$ . In order to validate the accuracy of the analysis, we also develop an event-driven simulation. We consider the following system configuration. The total channel number *C* of each cell is set to be 30, the voice call arrival rate  $\lambda_v$  can range from 0.04 to 0.29, the data call arrival rate  $\lambda_d$  can rage from 0.03 to 0.11. While other system parameters are fixed as:  $\mu_{vr} = 0.005$ ,  $\mu_{vh} = 0.01$ ,  $\mu_{dr} = 0.001$ ,  $\mu_{dh} = 0.0001$ .

## 4.1 Performance under different $\lambda_v$

Figs. 2 to 4 are the results under  $K_1 = 25$ ,  $K_2 = 25$  and  $\lambda_d = 0.03$ . Figs. 2 and 3 show that voice call dropping and blocking probability increase with the increase of data buffer size. This is because that under the buffer case, we place the data calls that cannot be serviced into the data buffer instead of simply blocking them. The data calls queuing in the data buffer will enter the system for service immediately if there is a channel becoming available. This means potentially more channels are occupied by the data calls. As the buffer size increases, more data calls can be

queued in the data buffer, and competes for channels with the voice call, thus the voice handoff dropping and blocking probability will be increased.

Observing from Fig. 4, it is clear that the data call blocking probability decreases with the increase of data buffer size. Under no buffer case, the data call will be blocked if there is no free channel available for data traffic, but in buffer case, such kinds of data calls can be buffered to wait for service, this result in the decrease of the data blocking probability. As the buffer size increases, more data calls can be queued in the buffer and less data calls will be blocked.

### 4.2 Performance under different $\lambda_d$

Figs. 5 and 6 presents the results obtained under  $\lambda_{vn} = \lambda_{vh} = 0.09$  and  $\mu_d = 0.006$ . Fig. 6 shows that the voice dropping probabilities may increase with the increase of  $\lambda_d$  and data buffer size. The increase of data buffer size means more data calls can be buffered, and obtain service if there are channels available. The results obtained also reveal that the voice call dropping will not be increase after the buffer size increases to a certain value, such as 10 in this case. This is caused by the traffic control policy itself, where only when the channels occupied are less than  $K_2$ , data call can be accepted. Therefore not every data call queued in the buffer can obtain service. This is crucial for maintaining the performance of handoff dropping probabilities. Similar results can be obtained for voice blocking probability under these conditions.

As expected, data blocking probability will decrease with the increasing of buffer size and  $\lambda_d$ , this is shown in Fig. 6.

### 4.3 Performance under different $K_2$

Figs. 7 and 8 indicate the system performance under different  $K_2$ . In the DTR-Q scheme,  $K_2$  is used to control the data traffic flow into the network. Fig. 7 shows that the voice dropping probabilities are increased with the increasing of  $K_2$  and data buffer size. In case that  $K_2$  is small when comparing with the channel capacity, such as  $K_2 = 8$ , the increment of voice dropping probability is not obvious with data buffer increase. This is because that more data calls will be blocked from entering the system if we set  $K_2$  to be small. But in case that  $K_2$  is large, the voice dropping probability will be increased significantly with the increasing data buffer size. This is caused by the fact that larger  $K_2$  means more data traffic can be accepted by the network system, the chance for data calls buffered in the data buffer to obtain service increase. Similar results can be obtained for voice blocking probability under these conditions.

Fig. 8 shows that the data call blocking probability is decreased with the increasing of data buffer size, the increase of data buffer size apparently benefits the data call and results in the decrease of data call blocking probability.

The increase of  $K_2$  causes more channels to be occupied by data and as a result, the data blocking probability is decreased.

## 5. CONCLUSION

In this paper, we propose a new handoff control policy for voice/data integrated wireless cellular networks, namely *Dual Trunk Reservation with Queuing* (DTR-Q). The performances of DTR-Q are analyzed with three-dimensional Markov process, and an event-driven simulation is developed to validate the accuracy of our analyzing. From the results obtained from both analysis and simulation, we can conclude that: 1) DTR-Q scheme can guarantee different *QoS* requirements from different traffic types; 2) buffering data call requests may benefit data users while impact negatively the performance of voice users.

There are a number of issues that will be addressed in our future research: 1) the bandwidth allocation schemes discussed in this paper are static. However as shown in [3, 8], dynamic schemes can always achieve better performance; 2) the handoff rates for both voice and data following convention have been assumed given, as there lacks adequate model that can derive handoff rates based on other parameters such as new call arrival rates and mobility, similar to those in [7]. Finally, we will be considering more realistic data arrival using TCP traffic, and exploring the self-similar behavior of the data arrival observed in [12].

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Fig. 1. The DTR-Q control scheme



Fig. 2. The voice dropping probability vs.  $\lambda_{\nu}$ 



Fig. 3. The voice blocking probability vs.  $\lambda_v$ 



Fig. 4. The data blocking probability vs.  $\lambda_v$ 



Fig. 5. The voice dropping probability vs.  $\lambda_d$ 



Fig. 6. The data blocking probability vs.  $\lambda_d$ 



Fig. 7. The voice dropping probability vs.  $K_2$ 



Fig. 8. The data blocking probability vs.  $K_2$