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Securing Physical and network layer using SNAuth-SPMAODV with DSSS for Mobile adhoc networks in Military Scenario

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Abstract-A mobile ad hoc network is an infrastructure less network, fast emerging today for deployment in variety of applications. During deployment, security emerges as a central requirement due to many attacks that affects the performance of the ad hoc networks. Particularly Denial of Service attack is one such severe attack against network and physical layer which is a challenging one to defend against in military communication environments.. The physical layer protocol in MANETs is responsible for bit-level transmission between network nodes and network layer is responsible to provide security services for both routing information and data message. This paper consider military scenarios and evaluate the performance of Security-enhanced-Multipath AODV (Ad hoc On-demand Distance Vector Routing) routing protocol called SNAuth-SPMAODV (Secure Neighbor Authentication Strict Priority Multipath Ad hoc On-demand Distance Vector Routing) with spread spectrum technology Direct Sequence Spread Spectrum (DSSS) to defend against signal jamming denial-of-service attacks in physical layer and network layer for MANET. The protocol discovers multiple paths between sender and receiver nodes without introducing extra packets into the network and authenticates the neighbor offering robustness in a secured MANET. SNAuth-SPMAODV with DSSS is found to be a good security solution even with its known security problems. The simulation is done using network simulator Qualnet 5.0 for different number of mobile nodes. The proposed model has shown improved results in terms of Average throughput, Average end to end delay, Average packet delivery ratio, Routing overhead and Average jitter.

Keywords- Mobile adhoc network, Denial of Service attack, Strict priority algorithn, Secure neighbor authentication, Direct Sequence Spread Spectrum

1. INTRODUCTION

In recent years, Mobile ad hoc Networks has started gaining attention from the industrial and academic research community due to their wide deployment and inherent nature of solving practical real world applications[5]. Many military and commercial applications have emerged due to the simplicity of the networks and widespread adoption of the technology. Most of the previous ad hoc network researchers have focused on problems such as routing and reliable communication, made a trusted environment. However, many applications in reality run only in untrusted environments and secured routing became a challenging one. Applications that may require secure communications include emergency response operations, military or police networks, safety critical business operations such as oil drilling platforms or mining operations [10]. For example, in emergency response operations such as the one after a natural disaster like a flood, tornado, hurricane and earthquake, when regular communication networks are damaged due to natural disasters, emergency rescue teams have to rely upon ad hoc networks for communication. To fend off malicious attackers in these emergency situations, many safety critical applications require secure communication. Ad hoc networks generally use a wireless radio communication channel. The main advantage of such networks is low cost deployment and maintenance. Today, the nodes and wireless hardware are inexpensive and readily available [11]. The network is automatically self configuring and self maintaining nature. Generally, wireless networks are vulnerable to several attacks. Particularly Denial of Service attack is one such severe attack against network and physical layer which is a challenging one to defend against. The physical layer protocol in MANETs is responsible for bit-level transmission between network nodes and network layer is responsible to provide security services for both routing information and data message [3]. The proposed model combines SNAuth-SPMAODV routing protocol with spread spectrum technology Direct Sequence Spread Spectrum (DSSS) to defend against signal jamming denial-of-service attacks in physical layer and network layer for MANET.



ISSN 2277-1956/V1N3-1840-1849

ISSN- 2277-1956

The paper is organized in such a way that Chapter 2 discusses Review of literature Chapter 3 discusses the proposed method, Chapter 4 discusses problem statement Chapter5 discusses simulation model and Chapter 6 gives experimental results, Chapter 7 discusses the conclusion

II. REVIEW OF LITERATURE

This chapter briefly describes the Denial of Service attacks for MANET.

A. Denial of Service attack

In this type of attack, an attacker attempts to prevent legitimate and authorized users from the services offered by the network. A denial of service (DoS) attack can be carried out in many ways. The classic way is to flood packets to any centralized resource present in the network so that the resource is no longer available to nodes in the network, as a result of which the network no longer operate in the manner in which it is designed to operate. This may lead to a failure in the delivery of guaranteed services to the end users. Due to the unique characteristics of ad hoc wireless networks, there exist many more ways to launch a DoS attack in such a network, which would not be possible in wired networks. DoS attacks can be launched against any layer in the network protocol stack. On the physical and MAC layers, an adversary could employ jamming signals which disrupt the on-going transmissions on the wireless channel. On the network layer, an adversary could take part in the routing process and exploit the routing protocol to disrupt the normal functioning of the network. For example, an adversary node could participate in a session but simply drop a certain number of packets, which may lead to degradation in the QoS being offered by the network. On the higher layers, an adversary could bring down critical services such as the key management service. For example, consider the following: In figure 1 assume a shortest path that exists from S to X and C and X cannot hear each other, that nodes **B** and **C** cannot hear each other, and that **M** is a malicious node attempting a denial of service attack. Suppose S wishes to communicate with X and that S has an unexpired route to X in its route cache. S transmits a data packet towards X with the source route S --> A --> B --> C --> D --> X contained in the packet's header. When M receives the packet, it can alter the source route in the packet's header, such as deleting D from the source route. Consequently, when C receives the altered packet, it attempts to forward the packet to X. Since X cannot hear C, the transmission is unsuccessful [6].

$S \leftrightarrow A \leftrightarrow B \leftrightarrow M \leftrightarrow C \leftrightarrow D \leftrightarrow X$

Figure 1: Denial of Service attack

B.Route Selection

Proactive routing protocols generate routes and store them for later use[7]. On- demand routing protocols only generate routes when necessary. The later is used more often in MANETs because they require fewer resources. The mostly used on-demand routing protocols are Ad-hoc On-demand Distance Vector (AODV) unless modified, the protocol use single routes between sender and receiver nodes. Multipath routing reduces dependency on single nodes and routes, offering robustness in a secured MANET[8].

C.Adhoc On demand Routing protocol (AODV)

AODV routing protocol is based on DSDV and DSR algorithm and is a state-of-the-art routing protocol that adopts a purely reactive strategy: it sets up a route on demand at the start of a communication session, and uses till it breaks, after which a new route setup is initiated [9]. This protocol is composed of two mechanism (1) Route Discovery and (2) Route Maintenance. AODV uses **Route Requ**est (RREQ), **Route Reply** (RREP) control messages in Route Discovery phase and **Route Error** (RERR) control message in Route Maintenance phase. The header information of this control messages can be seen in detail in [9]. In general, the nodes participating in the communication can be classified as source node, an intermediate node or a destination node. With each role, the behavior of a node actually varies. When a source node wants to connect to a destination node, first it checks in the existing route table, as to whether a fresh route to that destination is available or not. If a fresh enough route is available, it uses the same. Otherwise the node initiates a Route Discovery by broadcasting a RREQ control message to all of its neighbors. This RREQ message will further be forwarded (again broadcasted) by the intermediate nodes to their neighbors. This process will continue until the destination node or an intermediate node having a fresh route to the destination. At this stage eventually, a RREP control message is generated. Thus, a source node after sending a RREQ waits for RREPs to be received. Figure 2 depicts the traversal of control messages.



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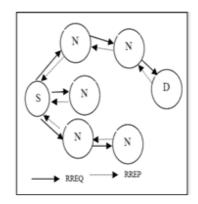


Figure 2: Traversal of Control Messages

D.Multipath Routing

Ad-hoc wireless routing protocols like AODV are mainly designed to discover and use a single route between a sender and receiver node. However, multiple paths between sender and receiver nodes can be used to offset the dynamic and unpredictable configuration of ad-hoc networks. They can also provide load balancing by spreading traffic along multiple routes, fault-tolerance by providing route resilience, and higher aggregate bandwidth.

Several multipath routing protocols based on DSR have been proposed, such as Split Multipath Routing (SMR) and Multipath Source Routing (MSR). Each of these multipath routing protocols broadcast data over all paths simultaneously. This technique has all the advantages previously mentioned, but it also introduces more packets into the MANET.

E. Strict-Priority Routing

Using multiple paths in ad-hoc networks to achieve higher bandwidth is not as straightforward as in wired networks. Because ad-hoc networks communicate over a wireless medium, radio interference may be a factor when a node communicating along one path interferes with a node communicating along another path, limiting the achievable throughput. Still, simulations have shown that broadcast multipath routing creates more overhead but provides better performance in congestion and capacity than unipath routing, provided the route length is within certain upper bound which is derivable. Additionally, the proper selection of routes using a strict priority multipath protocol can increase further the network throughput.

F.Secure Neighbor Authentication

The secure neighbor authentication has two variants. The first variant is based on *pair-wise shared secrets*, and the second variant is based on *certification*.

In secure neighbor authentication (SNAuth), every mobile node establishes an authenticated neighborhood on the move. Periodically, every mobile node X broadcasts its identity packet <SNAuth- HELLO, X> to its neighborhood.

1. In the pair-wise shared secret variant of SNAuth, Y, a neighboring receiver of the identity broadcast initiates a 3-way challenge-response handshake to authenticate X, the sender of the identity broadcast.

a. Suppose X and Y share a pair-wise secret k. Now Y selects a random nonce n1, encrypts n1 with k, sends the encrypted result ENC_k (n1) to X by a message <CHALLENGE, Y, ENC_k (n1)>.

b. If the receiver of the challenge message is indeed X, then it can decrypt ENC $_k$ (n1) and sees n1. X selects another random nonce n2, encrypts ENC $_k$ (n1 XOR n2), and sends back <RESPONSE1, X, n2, ENC $_k$ (n1 XOR n2)> as the response to the challenger Y.

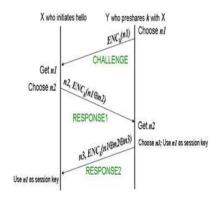
c. When Y receives the response, Y decrypts ENC_k (n1 XOR n2) and obtains n1 XOR n2. If Y can get the same result from XORing n2 in the response and its own challenge n1, then X passes the test with success. Otherwise, Y does not send any packet to X and does not receive packets from X except the response packets, until a correct <RESPONSE1> packet from X can pass the test. Upon detecting a success, Y puts X in its secure neighbor list. Y



selects a random nonce n3 and sends out a confirmation response <RESPONSE2, Y, n3, ENC_k (n1 XOR n2 XOR n3)> to X.

d. Upon receiving the RESPONSE2 message, X decrypts ENC_k (n1 XOR n2 XOR n3) and obtains n1 XOR n2 XOR n3. If this matches the result of XORing n1 that is previously decrypted, its own n2 and n3 in the RESPONSE2 packet, then X inserts Y into its secure neighbor list. (This three-way handshake is required because X needs to verify that Y actually knows k)

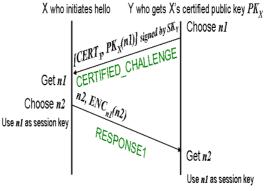
e. End of the challenge-response protocol. Figure 3 shows Challenge-Response Protocol-Three way handshake

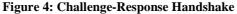




In the above description, all nonce length is currently set to 128-bit long. Encryption block length is 128-bit. Key k can be 128-bit, 192-bit, or 256-bit. Session key means that the key n1 is used until the time when the next HELLO received by Y from X successfully passes the test again.

2.A slightly different challenge-response scheme is used if Y does not pre-share a master secret k with X. Here X must broadcast its certificate $CERT_x = [X, certified public key PK_x, certificate valid time] in a CERTIFIED_HELLO message. For Y's CHALLENGE, Y uses PK_x to encrypt n1 and obtains ciphertext PK_x (n1). Y must also add its own certificate CERT_y = [Y, certified public key PK_y, certificate valid time] and sign the entire message with its own private key SKY. It recommend the public key cryptosystem in use be an Elliptic Curve Cryptosystem (ECC), because ECC features shorter certificate length and ciphertext length, thus incurring less communication overhead. Figure 4 shows Challenge-Response Handshake.$





When every neighboring receiver of X finishes the authentication and key-agreement process, node X obtains a secure snapshot of its neighborhood. In the neighborhood, every other node is authenticated and shares an IPsec



security association with the node X. As the SNAuth protocol runs on every mobile node, the statement is true if node X is replaced with any node X'.

G.Direct Sequence Spread Spectrum (DSSS) is a modulation technique. As with other spread spectrum technologies, the transmitted signal takes up more bandwidth than the information signal that is being modulated[3]. *Features*

DSSS phase-modulates a sine wave pseudo randomly with a continuous string of pseudo noise (PN) code symbols called "chips", each of which has a much shorter duration than an information bit. That is, each information bit is modulated by a sequence of much faster chips. Therefore, the chip rate is much higher than the information signal bit rate.

DSSS uses a signal structure in which the sequence of chips produced by the transmitter is known *a priori* by the receiver. The receiver can then use the same PN sequence to counteract the effect of the PN sequence on the received signal in order to reconstruct the information signal.

III. PROBLEM STATEMENT

This research investigates how to integrate security policies of a MANET with secure neighbor authentication that will allow the MANET to function securely in a military environment without degrading network performance. The specific problem to be addressed is how to use secure neighbor authentication of nodes in a multipath routing algorithm in MANET protected from Denial of service attack and physical layer security in military environment. Most of such performance analysis are normally done on commercial settings. For instance, wireless LAN technologies in the 2.4 GHz ISM frequency band are generally assumed, offering data rates up to 2 *Mbps* within the range of 250 *m*. This paper is motivated by the observation that such propagation and network models assumed by the current ad hoc networking simulations are quite different from real world military environments. In fact, a few hundred MHz frequency band (i.e., VHF or even HF) is used with very low data transmission rates (e.g., 384 *Kbps*) for the military scenarios. Table I summarizes these differences in terms of a physical layer model[12]. Networking environments such as network size, nodes' mobility model, and traffic patterns are quite different as well. For instance, the size of military networks is often far greater than that of their conventional counter parts both in the number of nodes and dimensions of the geographical areas.

| Table 1: physical layer model for military environments | | | | | |
|---|------------------|----------------------|--|--|--|
| Parameters | Military devices | Conventional devices | | | |
| Frequency | 30, 88, 300 MHz | 2.4, 5 GHz | | | |
| Propagation limits | -115 dBm | -110 dBm | | | |
| Radio propagation model | Two-ray ground | Line-of-sight | | | |
| Data rates | 9.6~384 Kbps | 2~54 <i>Mbps</i> | | | |
| Transmit power | 37 <i>dBm</i> | 15 <i>dBm</i> | | | |
| Receive sensitivity | -100 dBm | -90 dBm | | | |

Table I: physical layer model for military environments

IV. PROPOSED METHODOLOGY

A MANET is a collection of mobile routers that move dynamically in unpredictable directions. The links connecting the nodes are wireless and thus are not as dependable as wired links. The links are also susceptible to capacity constraints. A MANET environment is characterized by numerous security threats because the wireless links are vulnerable to Denial of service attack. The proposed method provide physical layer security and it reduces dependency on single nodes and routes; it discovers multiple paths between sender and receiver nodes it has the advantages of a multipath protocol without introducing extra packets into the network and authenticates the neighbor offering robustness in a secured MANET. It can be used to offset the dynamic and unpredictable configuration of



ad-hoc networks. They can also provide load balancing by spreading traffic along multiple routes, fault-tolerance by providing route resilience, and higher aggregate bandwidth in military environment.

V. SIMULATION MODEL

Using the QualNet network simulator [14], comprehensive simulations are made to evaluate the protocol. Qualnet provides a scalable simulation environment for multi-hop wireless ad hoc networks, with various medium access control protocols such as CSMA and IEEE 802.11. Channel and physical layer settings are modified to apply more realistic military scenarios. Note that PRC-999K device is used as a reference model. 802.11 DCF and UDP protocols are used for MAC and a transport protocols, respectively. Also, CBR traffic is utilized in the study. As the TCP-based application protocols such as telnet or FTP show unstable performance in mobile wireless communication, it can not evaluate precise performance of routing protocol itself. CBR application model sends one packet per second, which represents relatively low traffic patterns in military environments. Each packet size is 512 *Bytes*. In military environments, operational network size is very large as compare to conventional case. Nodes in the simulation are assumed to move according to the "random way point" mobility model. Pause time is fixed to 20 seconds. The attackers are positioned around the center of the routing mesh in all experiments.

To evaluate the performance of proposed method by 4 measurements: Packet delivery radio, average end-to-end delay, routing overhead and Throughput.

Results and Analysis

In this set of simulations, analyze performance of SNAuth-SPMAODV when the network size varies from 100 nodes to 1400 nodes. The network sizes and the respective network areas are shown in Table2 (approximately a walking Speed of soldiers). The size and the area are selected such that the node density is approximately constant, to properly evaluate proposed method.

| Nodes | Area (m) |
|-------|-----------|
| 100 | 1400×1400 |
| 200 | 2000×2000 |
| 400 | 2800×2800 |
| 600 | 3500×3500 |
| 800 | 4000×4000 |
| 1000 | 4500x4500 |
| 1200 | 4900x4900 |
| 1400 | 5300x5300 |

Table 2: Network sizes and areas.

EXPERIMENTAL RESULTS

Packet Delivery Ratio

From the Figure 5, it is shown that the proposed scheme (SNAuth-SPMAODV with DSSS) gives better Packet Delivery Ratio compared to SNAuth-SPMAODV with varying network size.



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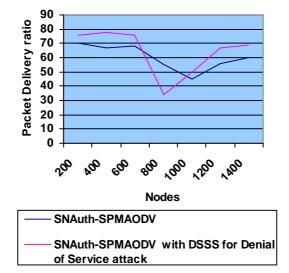


Figure 5 - SNAuth-SPMAODV with DSSS Packet delivery ratio

Throughput

Figure 6 demonstrates the throughput for SNAuth-SPMAODV with DSSS .It is clear that has a good performance compared to SNAuth-SPMAODV with DSSS with varying network size and malicious nodes.

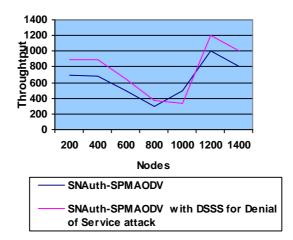
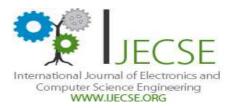


Figure 6 - SNAuth-SPMAODV with DSSS Throughput

Routing Overhead

Figure 7 show that Routing Overhead is lower in SNAuth-SPMAODV with DSSS for Denial of service attack and SNAuth-SPMAODV with varying network size and malicious nodes. Hence the Routing Overhead by malicious node has been minimized.



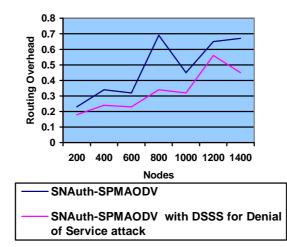


Figure 7- SNAuth-SPMAODV with DSSS Routing Overhead

8.4 Avg.End to End Delay

Figure 8 show that Avg.End to End Delay is lower in SNAuth-SPMAODV with DSSS for Denial of service attack and SNAuth-SPMAODV with varying network size and malicious nodes.

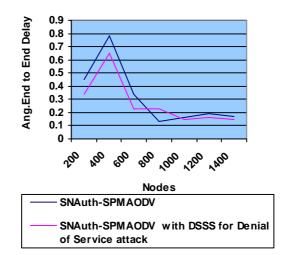
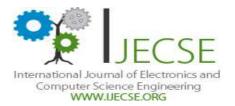


Figure 8- SNAuth-SPMAODV with DSSS Avg.End to End delay

Avg.Jitter

Figure 9 show that Avg.Jitter is lower in SNAuth-SPMAODV with DSSS for Denial of service attack and SNAuth-SPMAODV with varying network size and malicious nodes.



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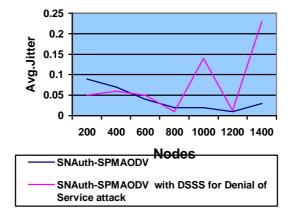


Figure 9- SNAuth-SPMAODV with DSSS Avg.jitter

CONCLUSION

Mobile Adhoc network is a collection of mobile nodes without infrastructure. During deployment, security emerges as a central requirement due to many attacks that affects the performance of the ad hoc networks. Particularly Denial of Service attack is one such severe attack against network and physical layer which is a challenging one to defend against. The physical layer protocol in MANETs is responsible for bit-level transmission between network nodes and network layer is responsible to provide security services for both routing information and data message. The proposed model combines SNAuth-SPMAODV routing protocol with spread spectrum technology Direct Sequence Spread Spectrum (DSSS) to defend against signal jamming denial-of-service attacks in physical layer and network layer for MANET. SNAuth-SPMAODV with DSSS is found to be a good security solution even with its known security problems.

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