

# Comments on “HEED : A Hybrid, Energy-Efficient, Distributed Clustering Approach for Ad Hoc Sensor Networks”

Chia-Hung Lin and Ming-Jer Tsai

Department of Computer Science  
National Tsing Hua University  
Hsing Chu, Taiwan, ROC  
E-mail: mjtsai@cs.nthu.edu.tw

## Abstract

*We provide a better sufficient condition for the connectivity of cluster heads asymptotically almost surely (a.a.s.) and a tighter bound on the number of cluster heads in HEED [1].*

**Index Terms**— Sensor networks, clustering, network lifetime, distributed algorithm.

## 1. Introduction

In network field  $[0, L]^2$ , where  $L \rightarrow \infty$ , Younis and Fahmy [1] show each  $\frac{R_c}{\sqrt{2}} \times \frac{R_c}{\sqrt{2}}$  cell

contains at least one node a.a.s. in Lemma 6. They claim that cluster heads are connected

a.a.s. in HEED if  $R_t \geq \sqrt{5} \left(2 + \frac{1}{\sqrt{2}}\right) R_c \approx 6R_c$  in Theorem 1. Besides, they show any

$\left(2 + \frac{1}{\sqrt{2}}\right) R_c \times \left(2 + \frac{1}{\sqrt{2}}\right) R_c$  area contains at least one cluster head a.a.s. in HEED in Lemma 7,

implying HEED has at least  $\frac{2}{4\sqrt{2}+9} \times \frac{L^2}{R_c^2} \approx 0.137 \frac{L^2}{R_c^2}$  cluster heads a.a.s. in network field

$[0, L]^2$ . In this paper, Theorem 1 shows cluster heads are connected a.a.s. in HEED if

$R_t \geq \left(1 + \frac{\sqrt{26+16\sqrt{2}}}{4}\right) R_c \approx 2.75R_c$  in Section 2, and Theorem 2 shows HEED has at least

$\frac{2}{3(4\sqrt{2}-1)} \times \frac{L^2}{R_c^2} \approx 0.144 \frac{L^2}{R_c^2}$  cluster heads a.a.s. in network field  $[0, L]^2$  in Section 3.

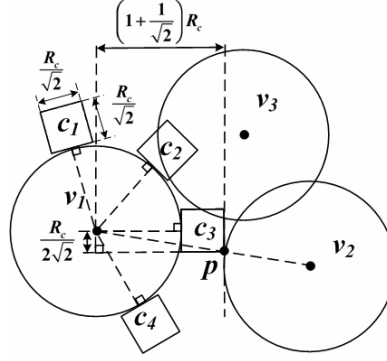


Fig. 1. Uncovered neighboring cells  $c_1$ ,  $c_2$ ,  $c_3$ , and  $c_4$  and a neighboring cluster head  $v_3$  of cluster head  $v_1$ , where the circle represents the coverage of the cluster head.  $v_2$  is a cluster head, where the straight line between  $v_1$  and  $v_2$  passes through  $c_3$ 's down-right corner  $P$  and the distance between  $P$  and  $v_2$  is  $R_c$ .

## 2. The Better Sufficient Condition

The following notations and lemmas are used in the proof of Theorem 1.

**Definition 1.** Let  $c$  be a  $\frac{R_c}{\sqrt{2}} \times \frac{R_c}{\sqrt{2}}$  cell and let  $v$  be a cluster head.  $c$  is said to be a  $v$ 's

uncovered neighboring cell if (1)  $c$  is not covered or partially covered<sup>i</sup> by  $v$ , and (2) the center of one  $c$ 's border is the point of tangency with  $v$ 's coverage.

**Definition 2.** Let  $v_1$  and  $v_2$  be two cluster heads.  $v_2$  is said to be  $v_1$ 's neighboring cluster head if  $v_2$  covers or partially covers at least one of  $v_1$ 's uncovered neighboring cells.

Take Fig. 1, for example.  $c_1$ ,  $c_2$ ,  $c_3$ , and  $c_4$  are some  $v_1$ 's uncovered neighboring cells, and  $v_3$  is  $v_1$ 's neighboring cluster head. Note that  $v_2$  is not  $v_1$ 's neighboring cluster head since  $v_2$ 's coverage does not cover or partially cover any  $v_1$ 's uncovered neighboring cell.

**Lemma 1.** Two neighboring cluster heads can communicate if  $R_t \geq \left(1 + \frac{\sqrt{26+16\sqrt{2}}}{4}\right) R_c$ .

**Proof.** The maximum distance between any two neighboring cluster heads is less than the

distance between  $v_1$  and  $v_2$  in Fig. 1, which is  $\left(1 + \sqrt{\left(\frac{1}{2\sqrt{2}}\right)^2 + \left(1 + \frac{1}{\sqrt{2}}\right)^2}\right) R_c$ . □

<sup>i</sup> A cell  $c$  is partially covered by  $v_1$  if some subarea of  $c$  is within  $v_1$ 's coverage.

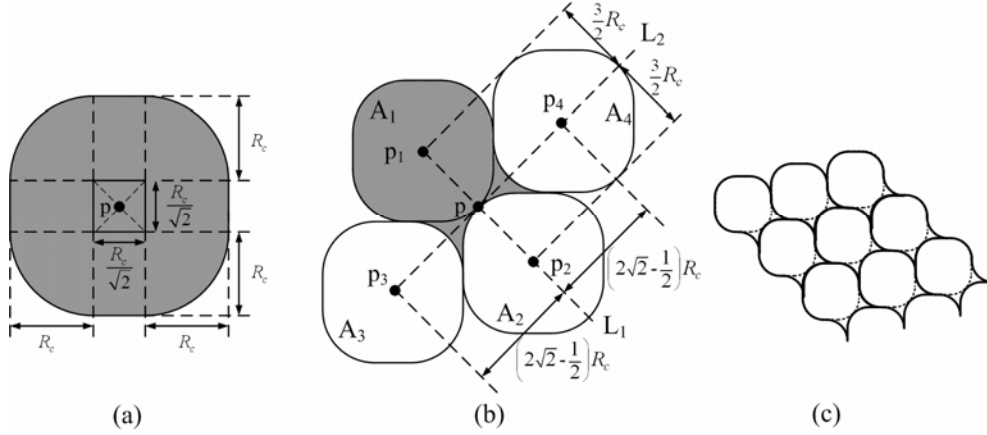


Fig. 2. Areas  $\theta$  and  $\Omega$ . (a) shows a  $\theta$  area shown in gray area. (b) shows a  $\Omega$  area shown in gray area. (c) shows the network field  $[0, L]^2$  divided into several  $\Omega$  areas.

**Lemma 2.** *Assume that HEED produces two isolated connected components of cluster heads  $G_1$  and  $G_2$  and that  $G_2$  lies on the right side of  $G_1$ . Then any cluster head  $v_1$  in  $G_1$  has one neighboring cluster head on its right side a.a.s.*

**Proof.** Let  $c$  be the rightmost one of  $v_1$ 's uncovered neighboring cells. So,  $c$  cannot be covered or partially covered by any cluster head on  $v_1$ 's left side. Since at least one node is in  $c$  a.a.s. by Lemma 6 in [1],  $c$  is covered or partially covered by at least one cluster head, say  $v_2$ , on  $v_1$ 's right side a.a.s. Clearly,  $v_2$  is  $v_1$ 's neighboring cluster head.  $\square$

**Theorem 1.** *HEED produces a connected multihop network a.a.s. if  $R_c \geq \left(1 + \frac{\sqrt{26+16\sqrt{2}}}{4}\right)R_c$ .*

The proof of Theorem 1 is similar to that of Theorem 1 in [1], and hence is omitted here.

### 3. The Tighter Bound

The following notation and lemma are used in the proof of Theorem 2.

**Definition 3.** *Area  $A$  is said to be a  $\theta$  area (as shown in Fig. 2a) if a  $\frac{R_c}{\sqrt{2}} \times \frac{R_c}{\sqrt{2}}$  cell  $c$  exists such that  $c$  is only covered or partially covered by any cluster head inside  $A$ , where  $c$  is said to be centrally enclosed by  $A$ . In addition, the intersection point  $p$  of two diagonals of  $c$  is said to be  $A$ 's core.*

**Definition 4.** Let  $p_1, p_2, p_3$  and  $p_4$  be the cores of four non-overlapping  $\theta$  areas  $A_1, A_2, A_3$  and  $A_4$ , respectively, where  $A_1$  tangent to  $A_2$  in one point  $p$ ,  $D(p_1, p_2) = 3R_c$  ( $D(u, v)$  denotes the distance between two points  $u$  and  $v$ ),  $D(p_3, p) = D(p_4, p) = \left(2\sqrt{2} - \frac{1}{2}\right)R_c$ ,  $p_1, p_2, p$  lie in line  $L_1$ ,  $p_3, p_4, p$  lie in line  $L_2$ , and  $L_1$  is perpendicular to  $L_2$ . Then an area is a  $\Omega$  area (as shown in Fig. 2b) if it contains  $A_1$ , the area surrounded by  $A_1, A_2, A_3$ , and the area surrounded by  $A_1, A_2, A_4$ .

**Lemma 3.** Any  $\theta$  area contains at least one cluster head a.a.s.

**Proof.** Let  $A$  be a  $\theta$  area centrally enclosing  $\frac{R_c}{\sqrt{2}} \times \frac{R_c}{\sqrt{2}}$  cell  $c$ . Since at least one node is in  $c$  a.a.s. by Lemma 6 in [1] and  $c$  cannot be covered or partially covered by any cluster head outside  $A$ , there is at least one cluster head inside  $A$  a.a.s.  $\square$

**Theorem 2.** In network field  $[0, L]^2$ , where  $L \rightarrow \infty$ , the number of cluster heads in HEED

is at least  $\frac{2}{3(4\sqrt{2}-1)} \times \frac{L^2}{R_c^2}$  a.a.s.

**Proof.** The network field  $[0, L]^2$  can be divided into several  $\Omega$  areas as shown in Fig. 2c.

Because a  $\Omega$  area contains a  $\theta$  area, any  $\Omega$  area contains at least one cluster head

a.a.s. by Lemma 3. The measurement of a  $\Omega$  area is  $\left(6\sqrt{2} - \frac{3}{2}\right)R_c^2$ . Thus there is at least

$\frac{2}{3(4\sqrt{2}-1)} \times \frac{L^2}{R_c^2}$  cluster heads a.a.s. in network field  $[0, L]^2$ .  $\square$

## References

- [1] O. Younis and S. Fahmy, "HEED: A Hybrid, Energy-Efficient, Distributed Clustering Approach for Ad Hoc Sensor Networks," *IEEE Trans. Mobile Computing*, vol. 3, no. 4, pp. 366-379, Oct.-Dec. 2004.